

Geometric design rules

Postlab lab 2

Prelab lab 3

Week 4

- Monday lab2
 - Carry gate schematic
- Tuesday
 - Lecture Layout of CMOS gates
- Thursday
 - Postlab review lab 2
 - Geometrical design rules incl. prelab 3
 - Tutorial POTW Layout (Victor)
- Friday Deadline prelab 3
 - Layout of carry ckt of full adder

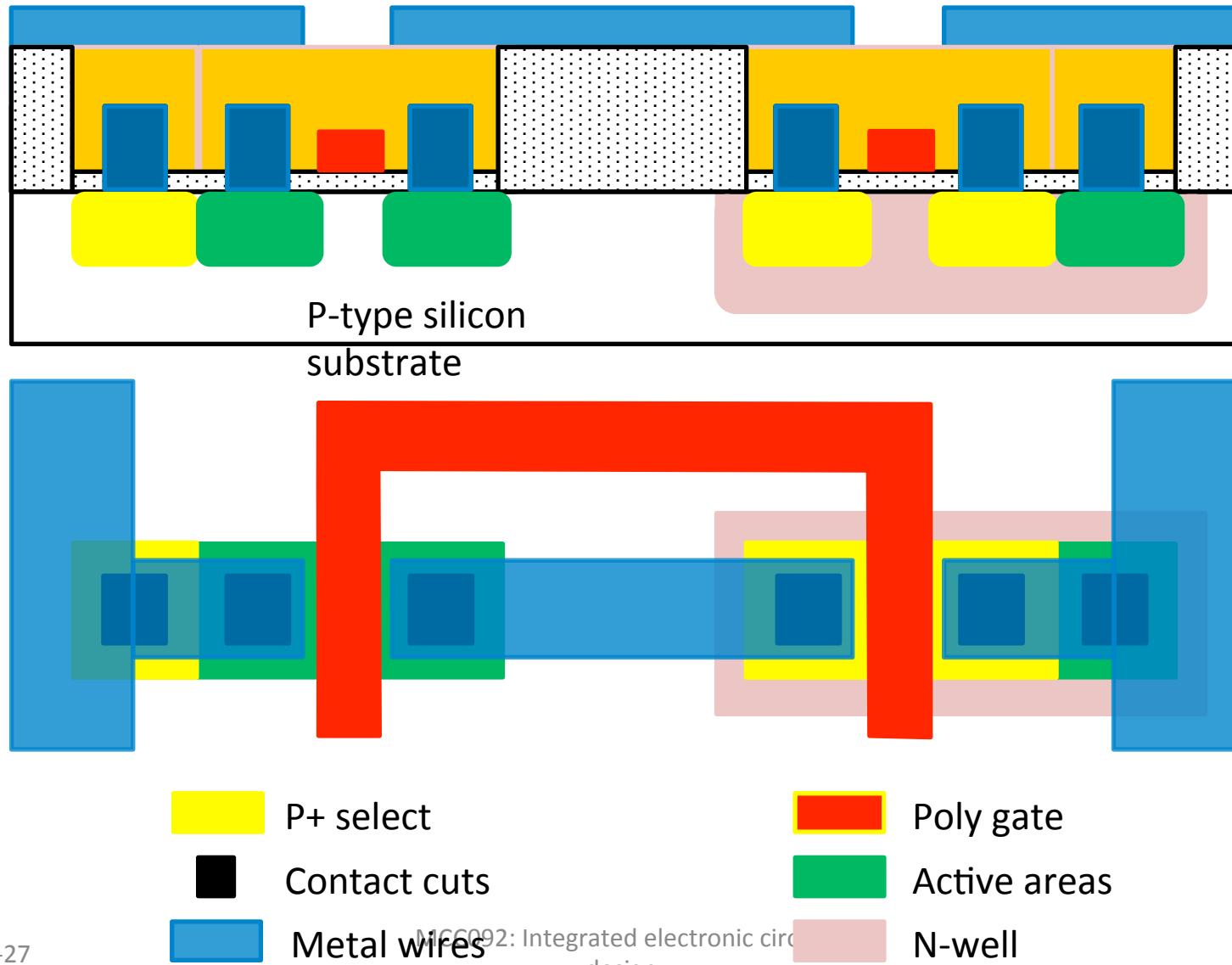
Practical things

- Thursday morning consultation time moved to 9.15-10 starting next Thursday
- Lars Svensson will help out with lab 3
- Mid-course meeting with student representatives held today.
 - Cadence remotely with VPN?
 - Chapter 2 problems?

From MUD cards

- How should metal line overlap contacts?
- Why do we connect input to V_{DD} on the second connection point on P-substrate?
- Euler path
 - How to draw it through odd and even nodes.
 - How to convert path to layout.

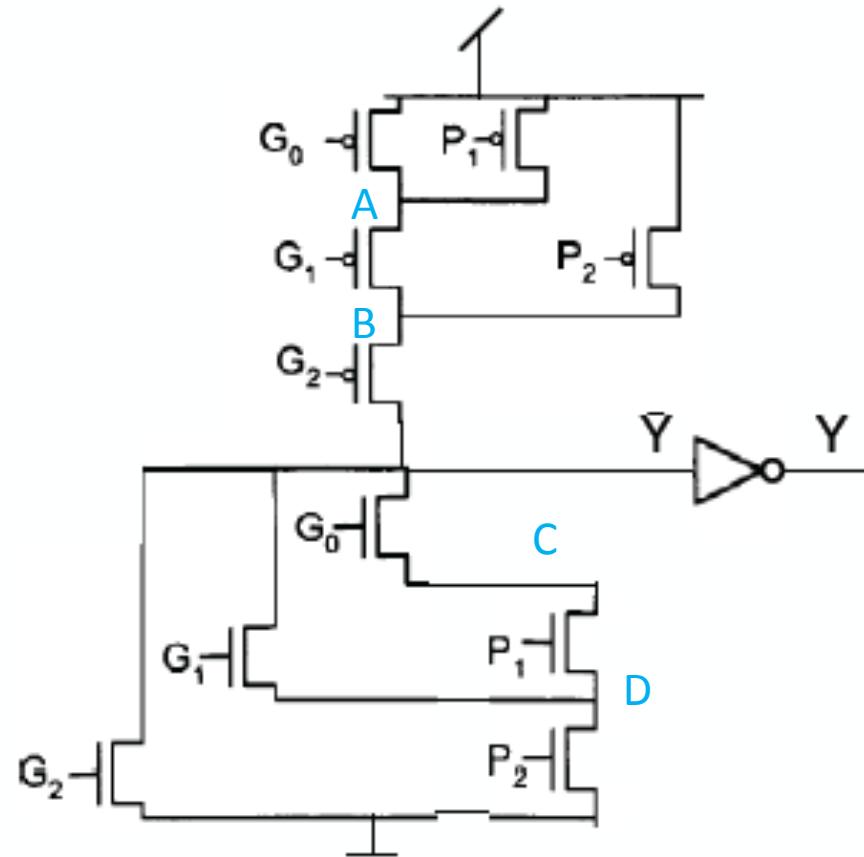
Inverter mask set and fabrication



Example Euler paths

It is a good idea to label the circuit nodes that are neither the output, nor VDD or VSS.

I have done so here.
There are four such nodes:
A, B, C, D.



Example Euler paths

How many transistors are connected to each of the circuit nodes in the **p-net**?

VDD: 3

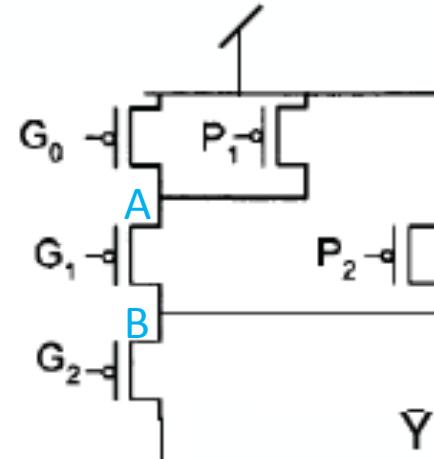
A: 3

B: 3

Y': 1

For an Euler path to exist there can only be 0 or 2 vertices (that is circuit nodes) with odd number of edges (that is transistors).

So with this schematic for the p-net there is no Euler path.



Example Euler paths

How many transistors are connected to each of the circuit nodes in the **n-net**?

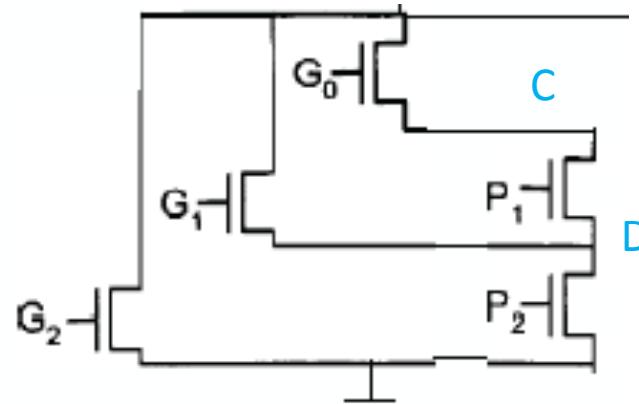
VSS/gnd: 2

D: 3

C: 2

Y': 3

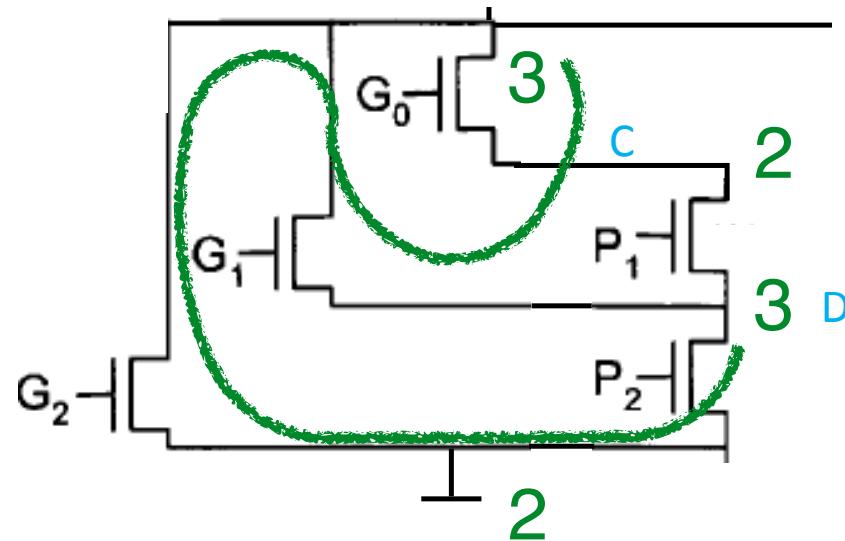
There are two vertices (circuit nodes) with an odd number of edges (transistors). Thus, an Euler path can be formed. It has to start and end in the circuit nodes with odd number of transistors: The output, Y, and internal node D.



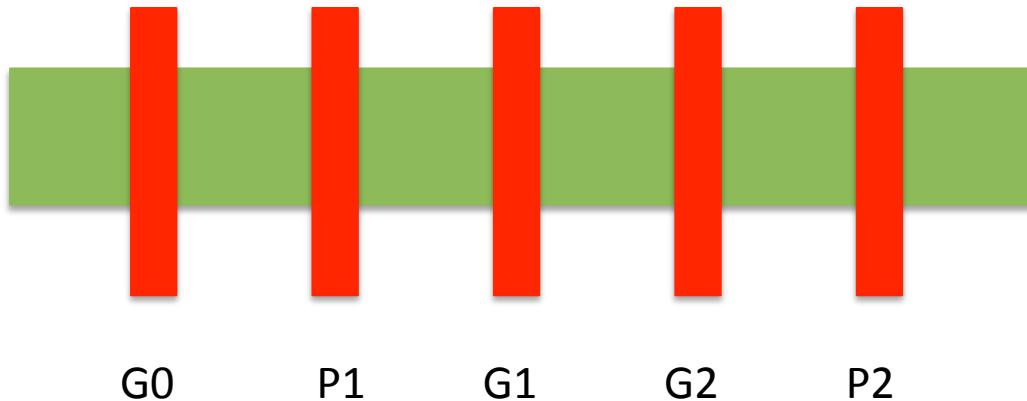
Example Euler paths

We have to start in a odd node and end in the other odd node.

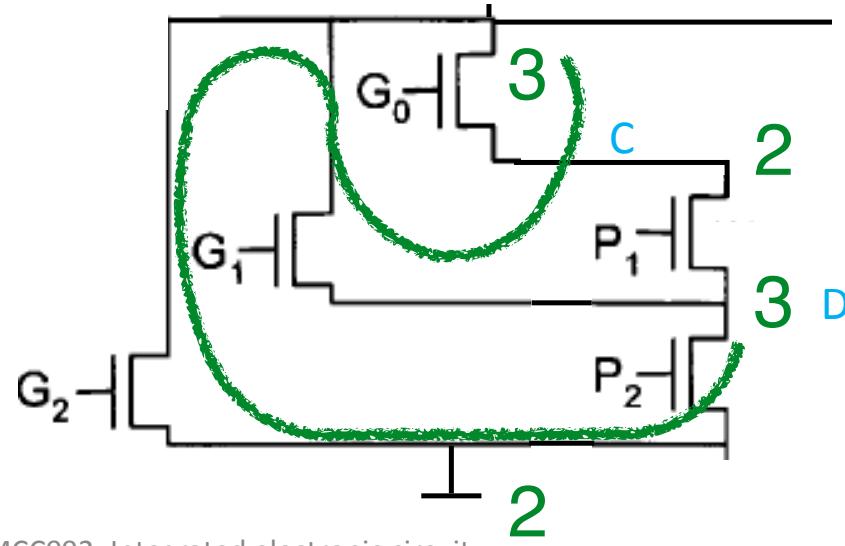
Here is one solution.
There are several
other possibilities.



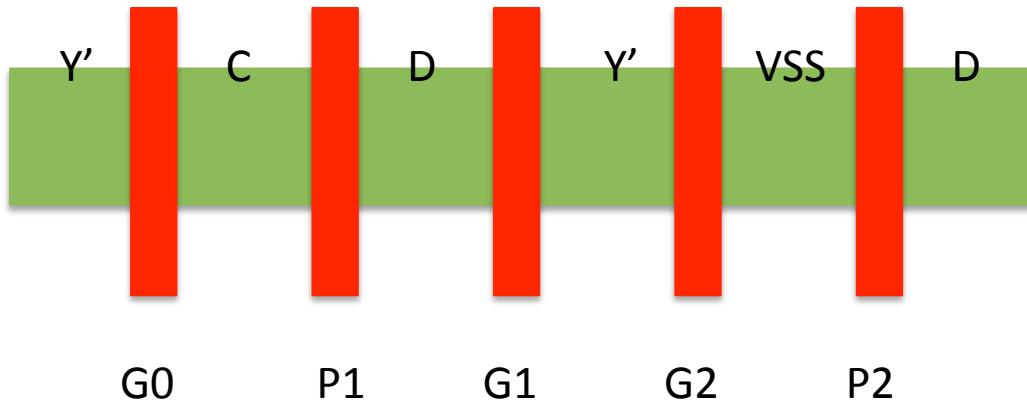
Example Euler paths



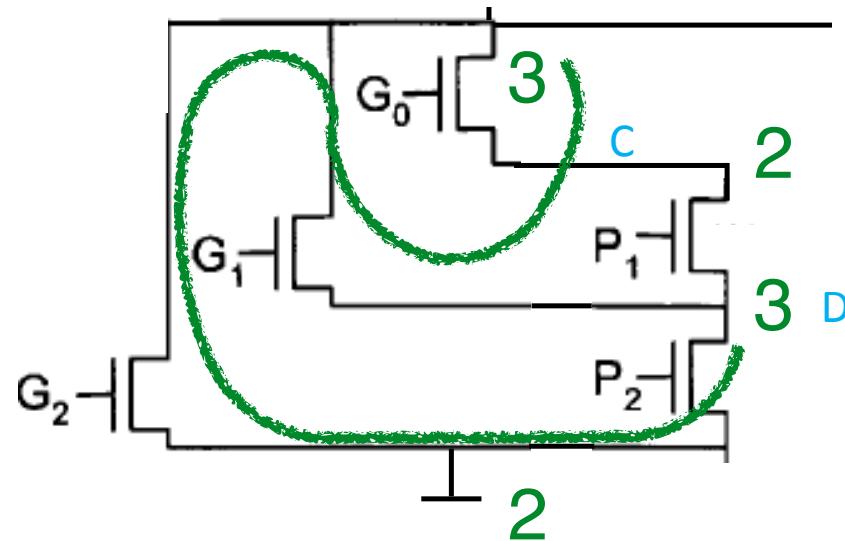
To map Euler path to a continuous line of diffusion start by labeling the transistors with their signal name starting at one end of Euler path.
(If some signals appear more than once give unique names)



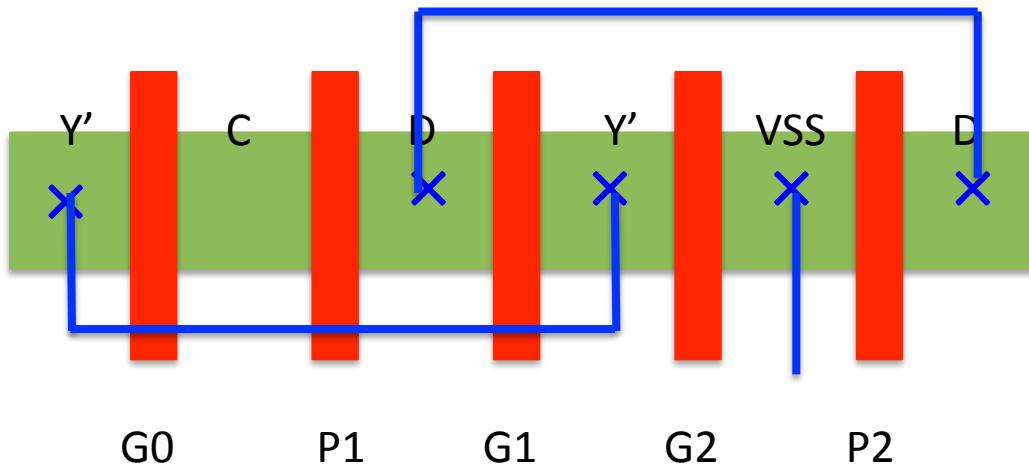
Example Euler paths



Then label also the circuit nodes along the path.

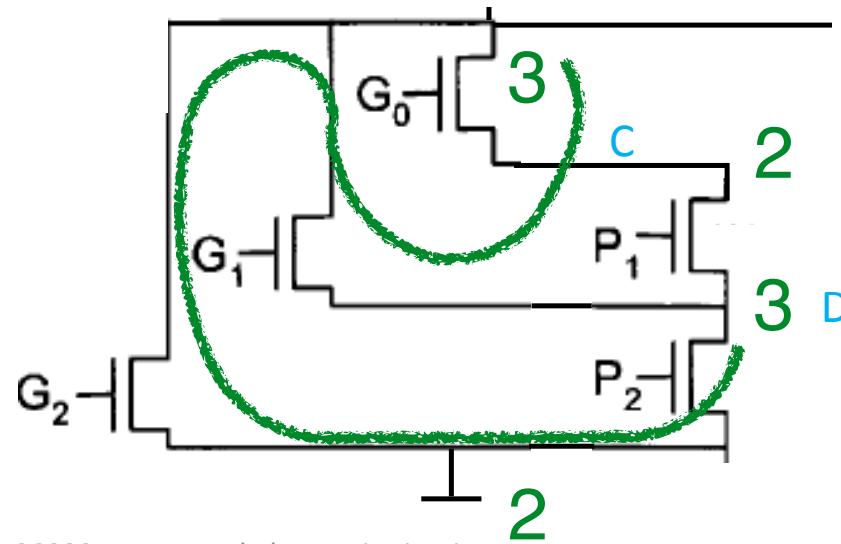


Example Euler paths



The make the connections for the source-drain areas that have the same circuit-node label.

I just draw it as a stick diagram here.



Example with quiz

Draw the schematic for a static CMOS gate with the logical function:

$$f = \overline{(A \cdot B + C \cdot D) \cdot (E + F + G)}$$

- 1 The n-net can be laid out with single line of diffusion. T/F
- 2 The p-net can be laid out with single line of diffusion. T/F
- 3 Can you make the order ABCDEFG work for the n-net?
- 4 Can you make the order ABCDEFG work for the p-net?

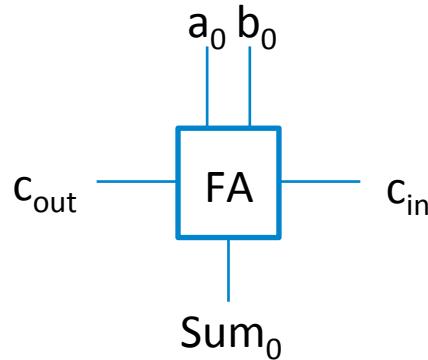
Postlab 2

Prelab 2

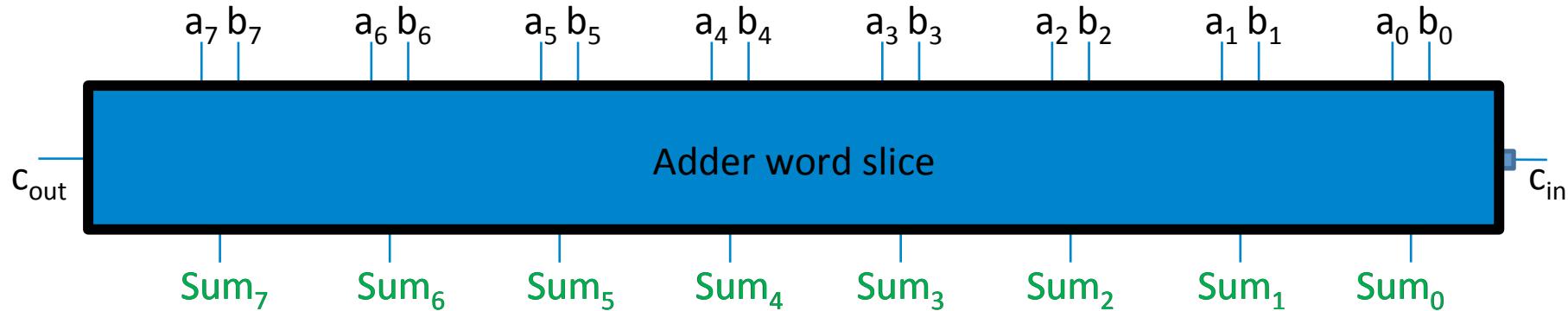
- Design task: design with MOSFETs the carry bit-cell for an 8-bit ripple-carry adder!
- The 8-bit carry logic is to be implemented by an iterative logic array consisting of eight instances of the bit-cell that you have designed.
- The carry bit-cell has three inputs (two bits a , b and a carry-in memory bit), and one output, the carry-out memory bit to the next more significant bit.

Iterative logic arrays: FULL ADDER

Find Boolean equations
from truth table!

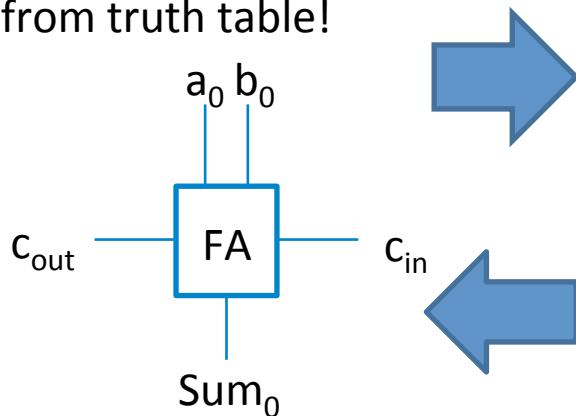


BOOLEAN TRUTHTABLE				
A	B	CIN	COUT	SUM
0	0	0	0	0
0	0	1	0	1
0	1	0	0	1
0	1	1	1	0
1	0	0	0	1
1	0	1	1	0
1	1	0	1	0
1	1	1	1	1



Iterative logic arrays: FULL ADDER

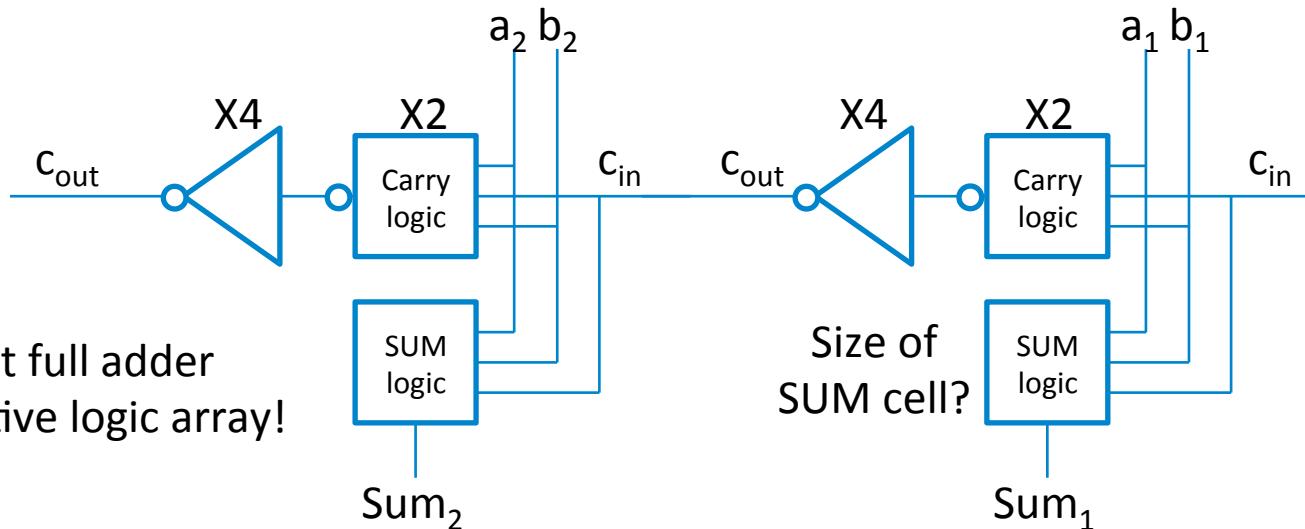
Find Boolean equations from truth table!



BOOLEAN TRUTH TABLE				
A	B	CIN	COUT	SUM
0	0	0	0	0
0	0	1	0	1
0	1	0	0	1
0	1	1	1	0
1	0	0	0	1
1	0	1	1	0
1	1	0	1	0
1	1	1	1	1

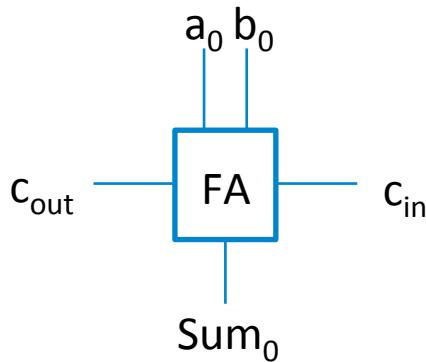
Design full adder by divide and conquer

Design 8-bit full adder as an iterative logic array!

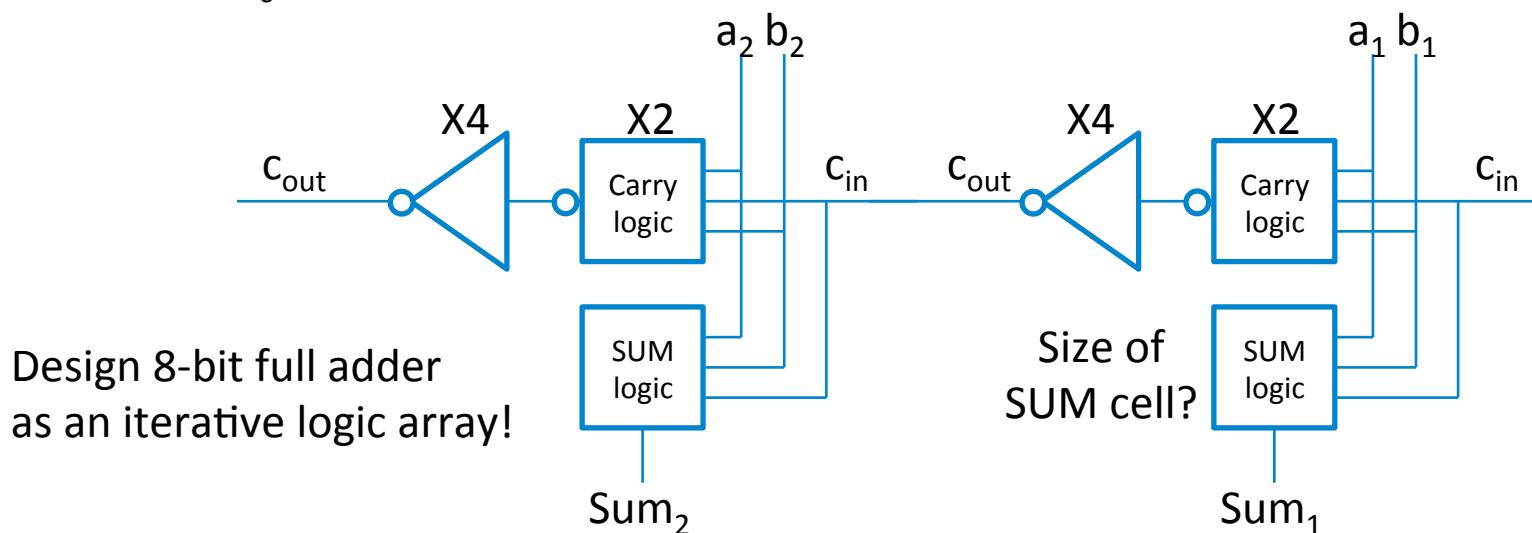


Iterative logic arrays: FULL ADDER

Find Boolean equations
from truth table!



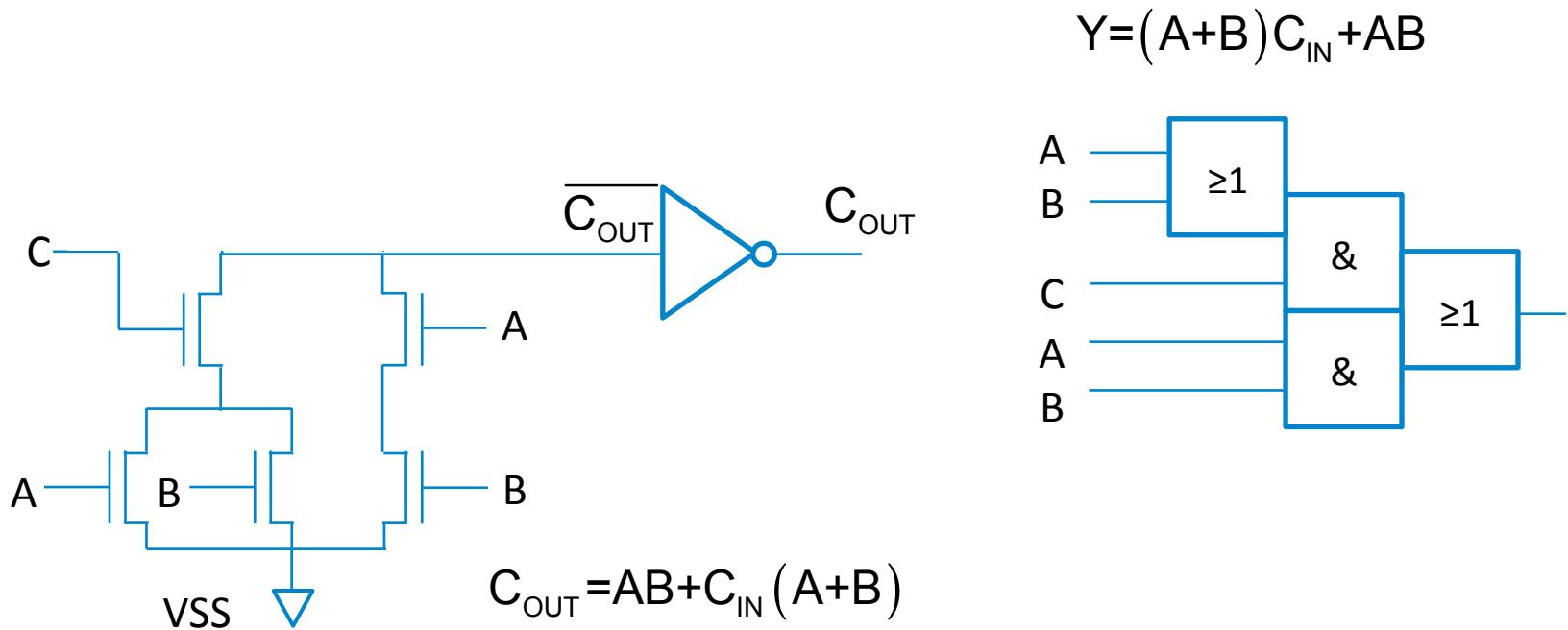
Propagation delay estimation:
Assume carry cell has logical effort g and
parasitic delay p .
8-bit adder delay: A first crude estimate!
 $d=8\times(p+g\times2+p_{inv}+1/2)\approx8\times10=80 \rightarrow t_{pd}=400 \text{ ps}$
Corresponding to 1.25 GHz clock frequency
But load from SUM cell was neglected and will
slow down the design!



Design 8-bit full adder
as an iterative logic array!

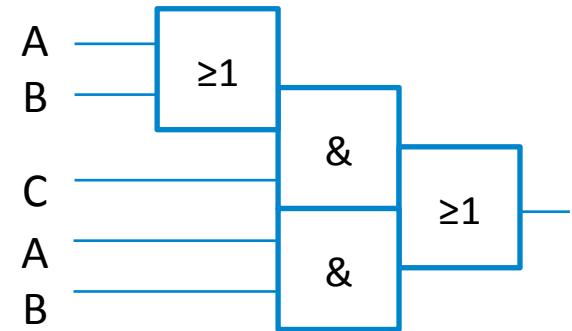
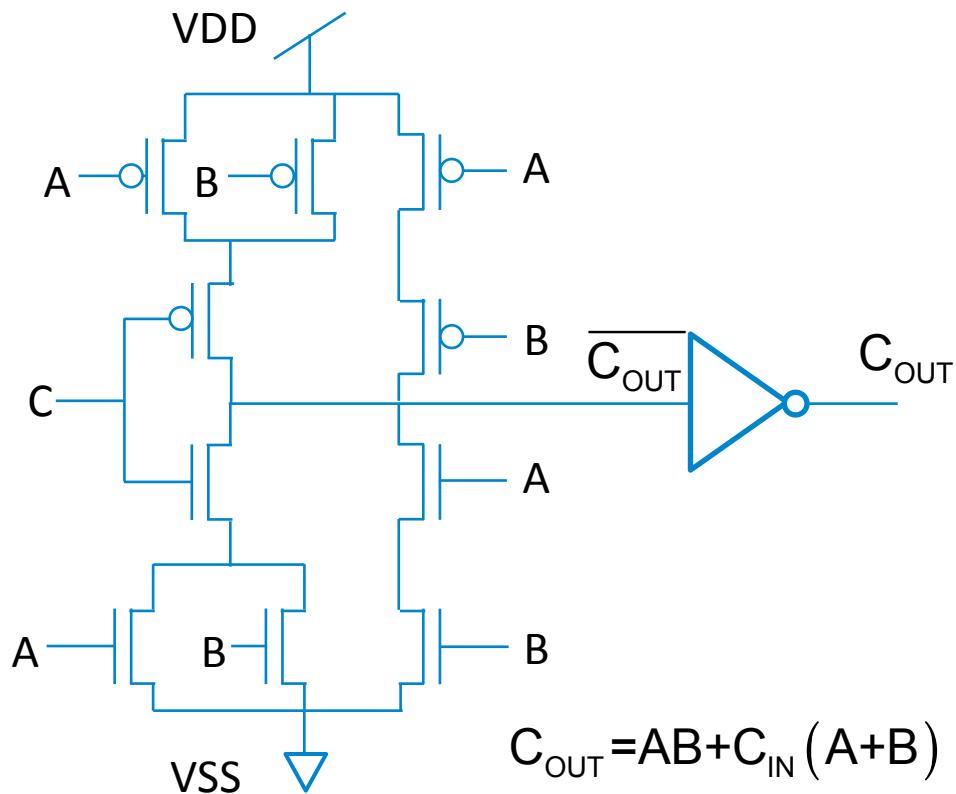
Designing the carry cell

$$\overline{C_{OUT}} = (\overline{A} + \overline{B})(\overline{C_{IN}} + \overline{AB}) = \overline{AB} + \overline{C_{IN}}(\overline{A} + \overline{B})$$



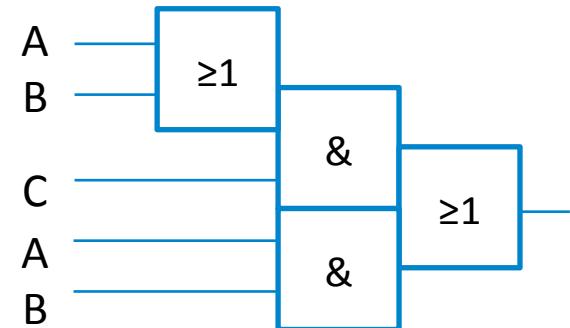
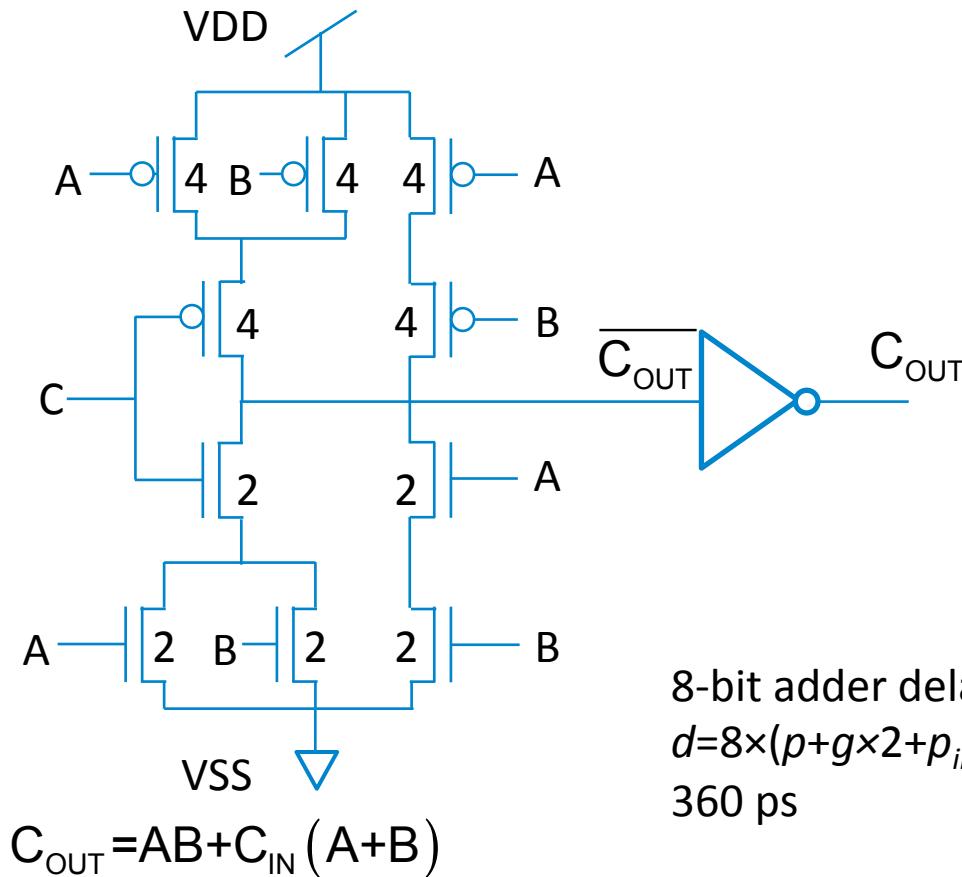
Designing the carry cell

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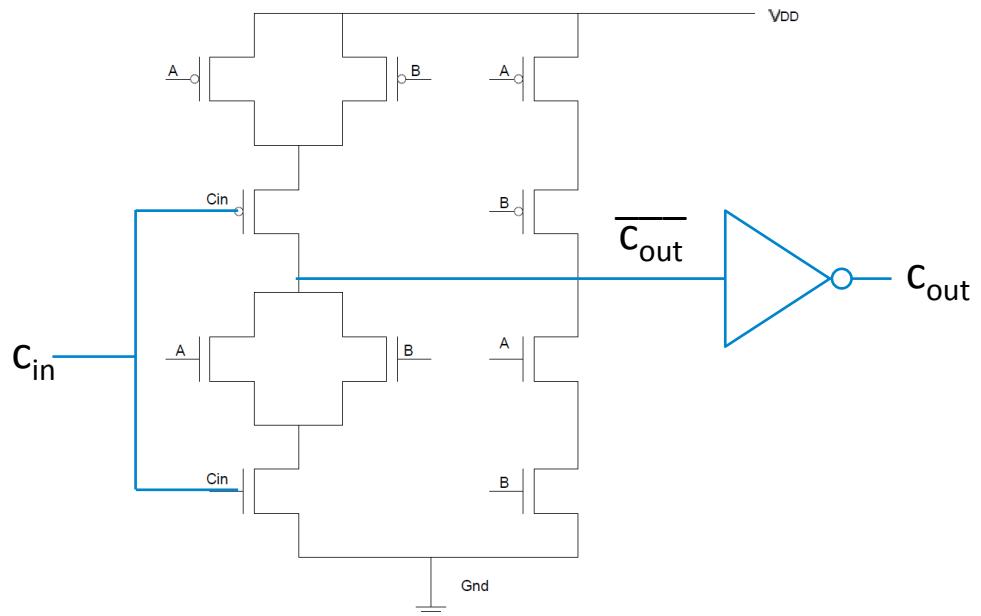
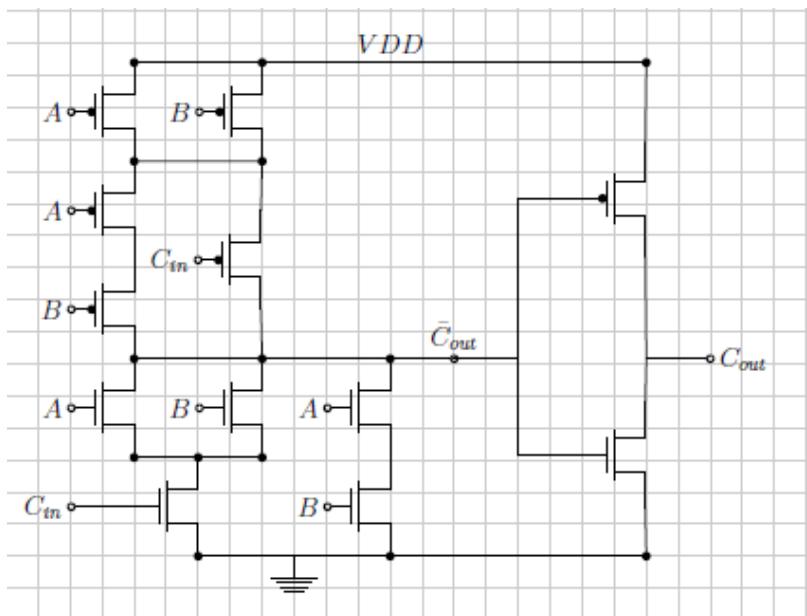
Designing the carry cell

$$\overline{C_{OUT}} = (\overline{A} + \overline{B})(\overline{C_{IN}} + \overline{AB}) = \overline{AB} + \overline{C_{IN}}(\overline{A} + \overline{B})$$

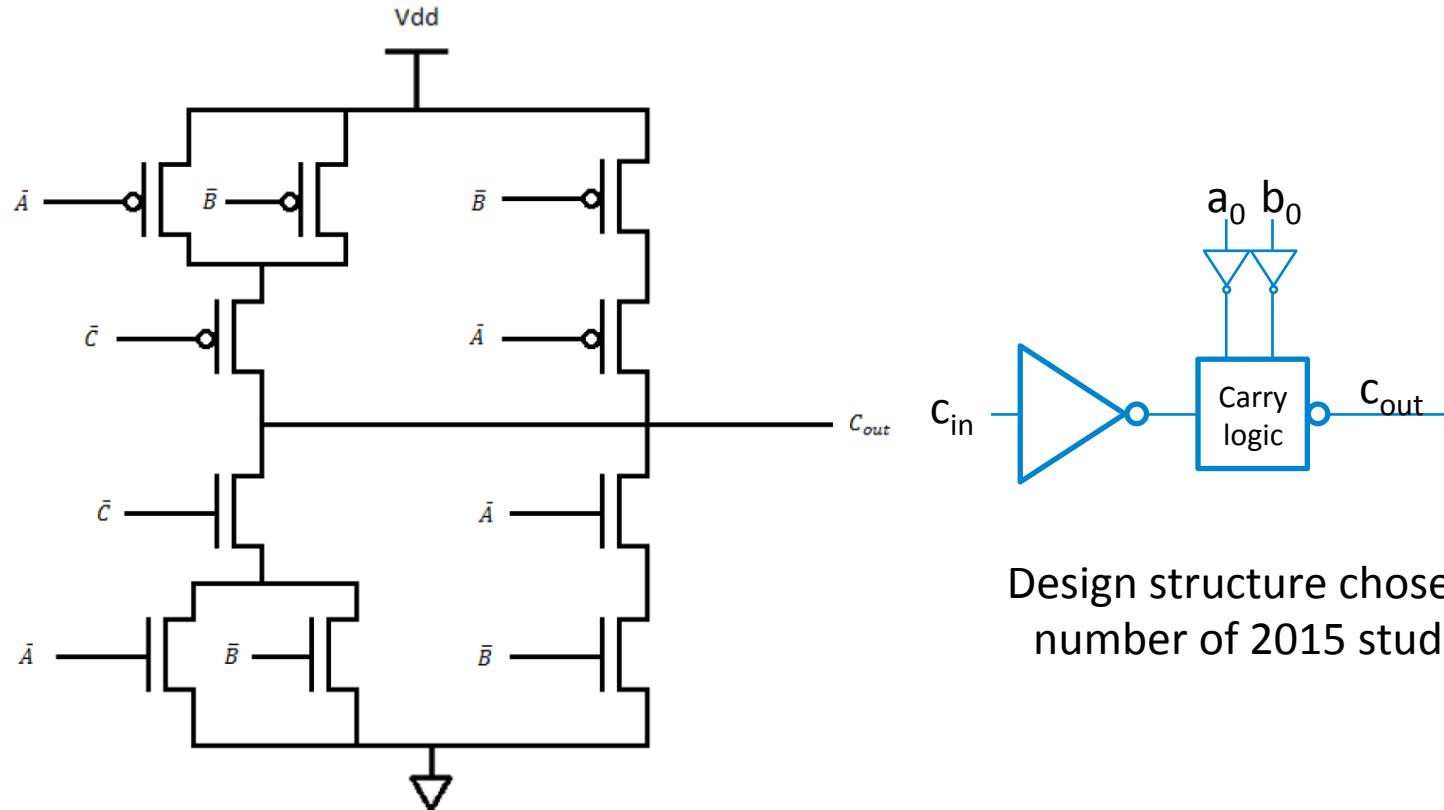


8-bit adder delay: Improved estimate w pinv = 0.9
 $d=8\times(p+g\times2+p_{inv}+1/2)\approx8\times(5p_{inv}+2\times2+0.5)=72\rightarrow t_{pd}=360\text{ ps}$

2015 hand-in examples

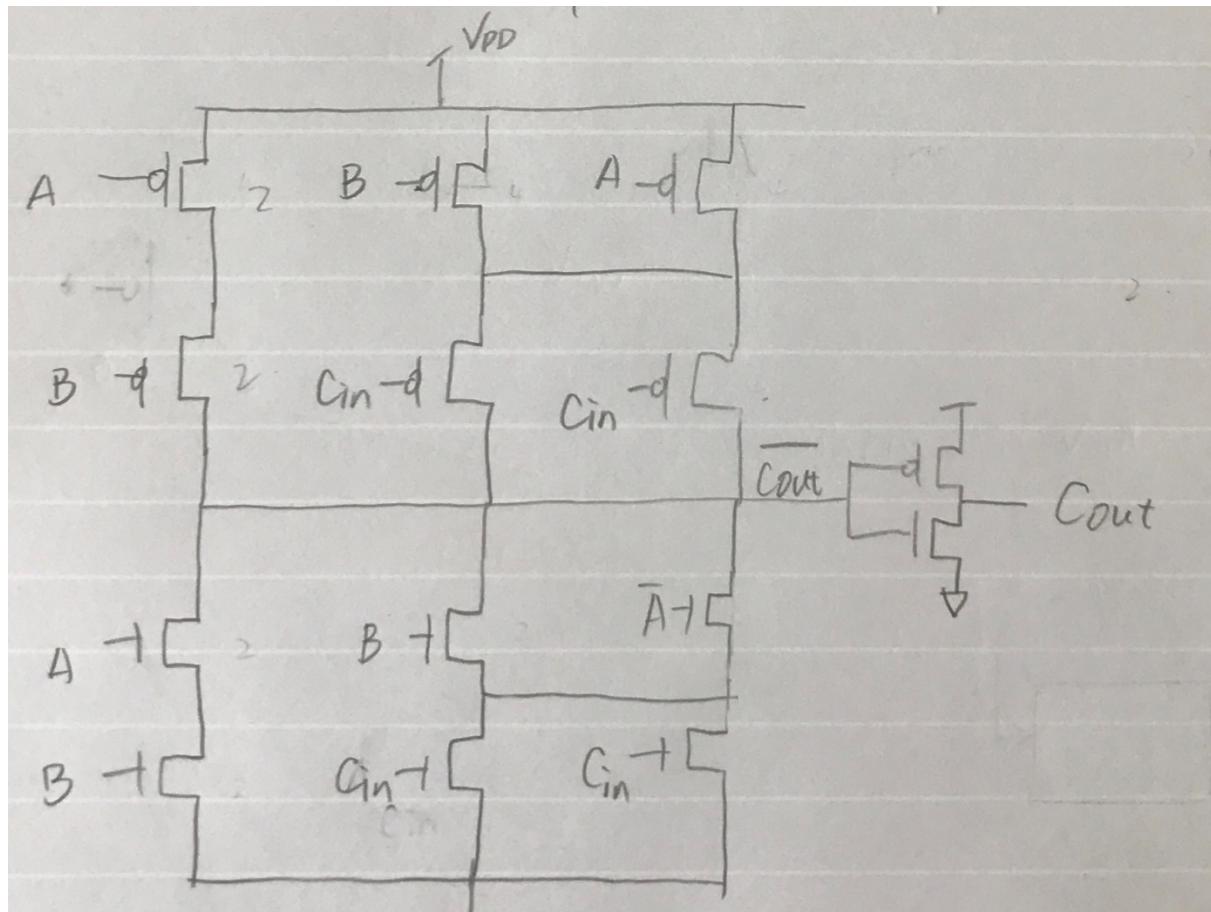


2015 hand-in examples

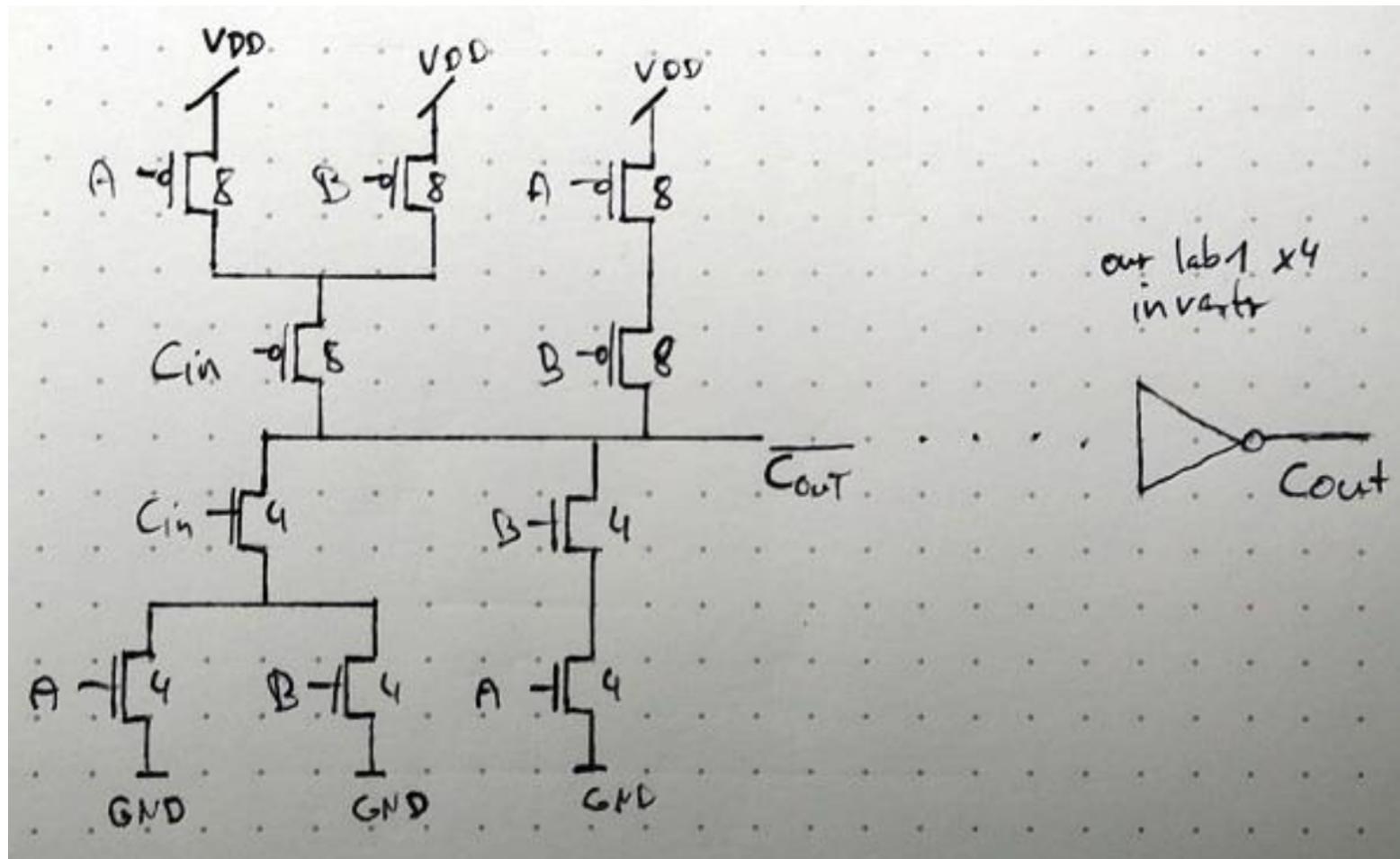


Design structure chosen by a
number of 2015 students!

2017 hand-in example



2017 hand-in example



Transient Response

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Transient Response

Name Vis

A

B

/CIn

/Cout



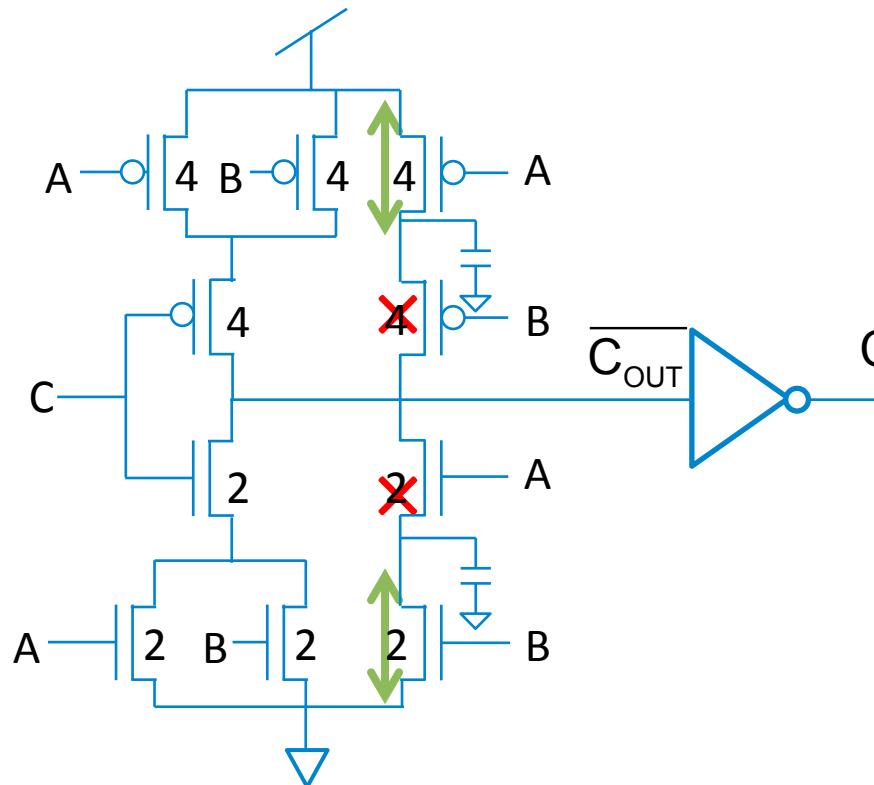
Simulating the carry cell

A and B inputs:

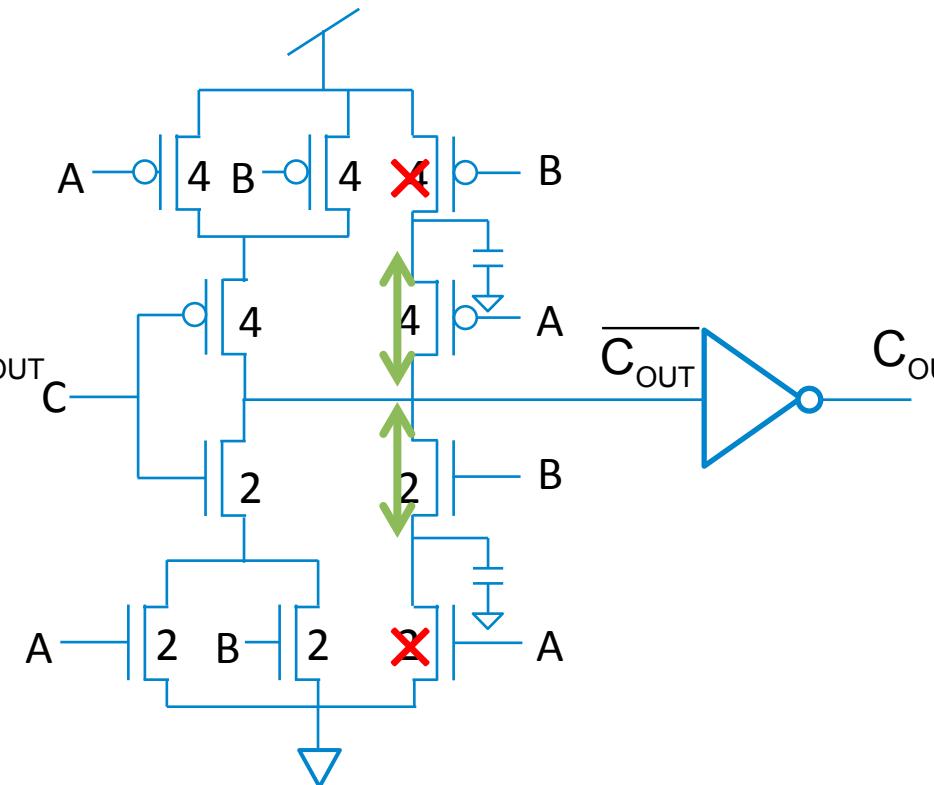
$B<0:7> = (1,1,1,1,1,1,1,1)$

$A<0:7> = (ln1,0,0,0,0,0,0,ln7)$

Impact of test signals



Expectation: fast



Expectation: slow

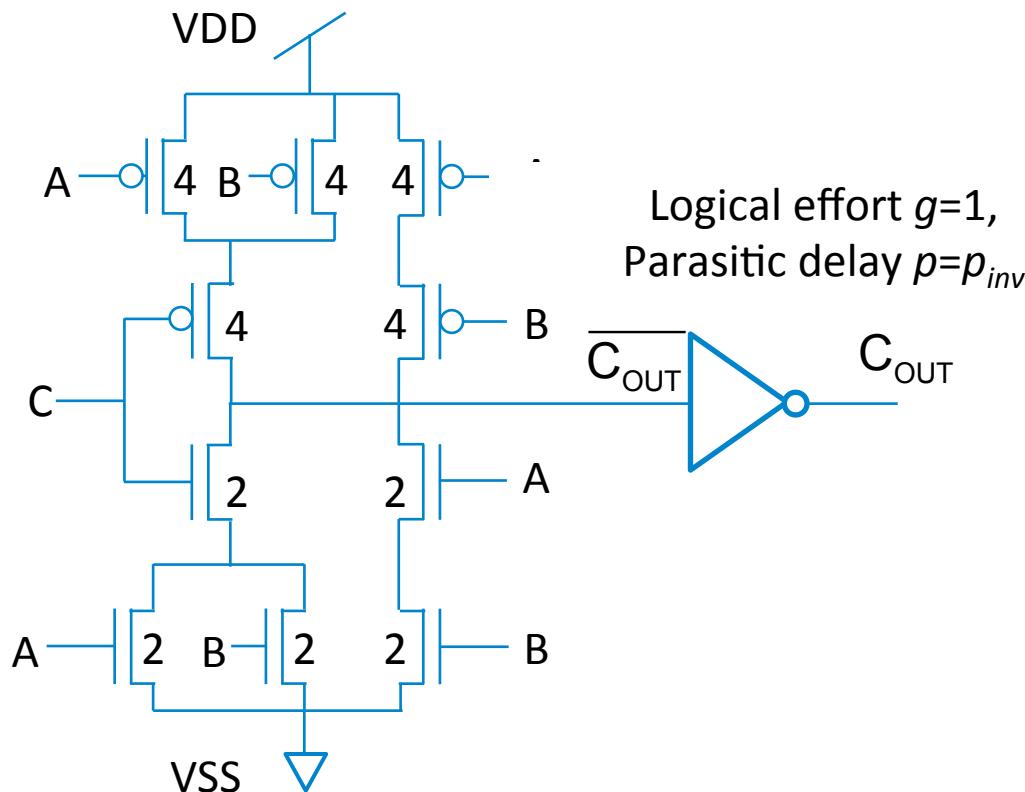
Simulated propagation delays TT models

Configuration	BAAB	ABBA	BABA	ABAB
Cin-Cout rise [ps]	430	393	453	371
Cin-Cout fall [ps]	446	453	488	413

Conclusion: Symmetrical schematics have less delay variation and the worst-case delay is shorter

Optimizing the size of the inverter

Logical effort $g=2$,
Parasitic delay $p=4p_{inv}$



$$C_{OUT} = AB + C_{IN}(A+B)$$

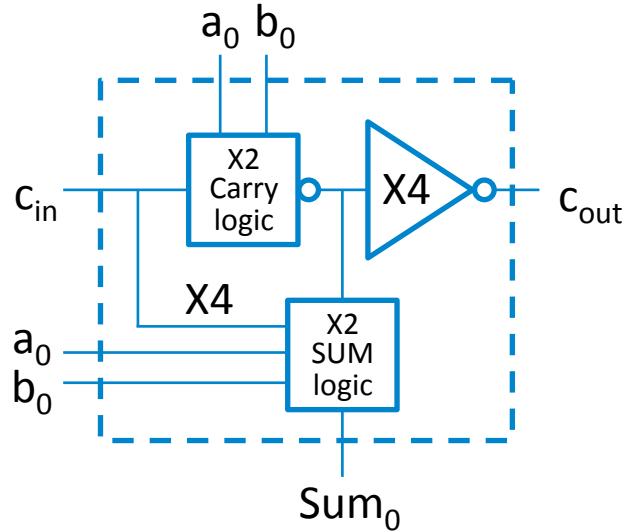
$$\begin{aligned}
 \text{Optimal path delay:} \\
 G &= 2 \times 1 = 2 \\
 H &= C_L/C_{IN} = C_{IN}/C_{IN} = 1 \\
 F &= 2 \times 1 = 2 \\
 f_{opt} &= \sqrt{2}
 \end{aligned}$$

Sizing: Gate $\sqrt{2}$ larger than inverter

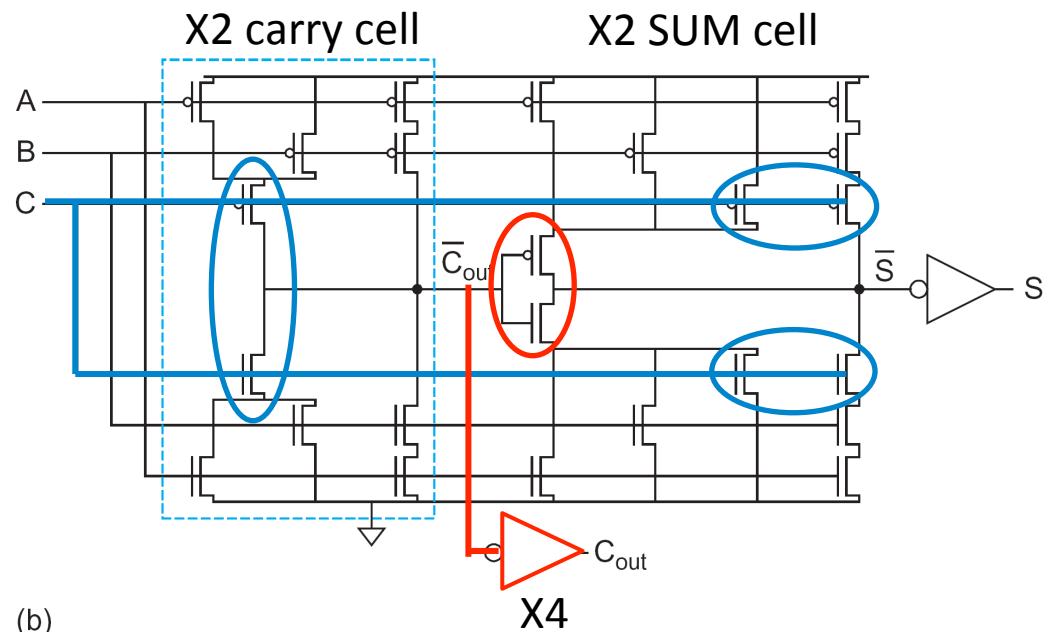
$$h_{\text{gate}} = 1/\sqrt{2}$$

$$h_{\text{inv}} = \sqrt{2}$$

Textbook solution



Design structure as proposed in
textbook!

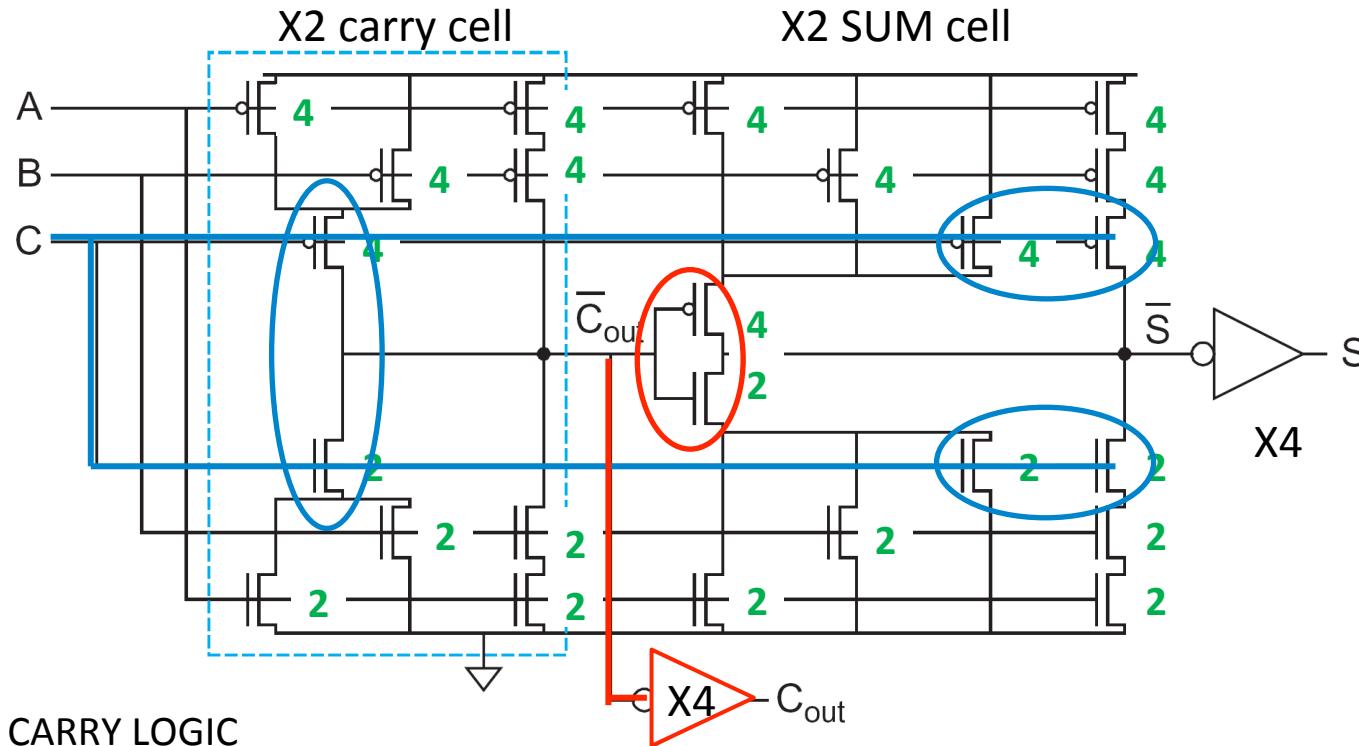


(b)

Textbook solution

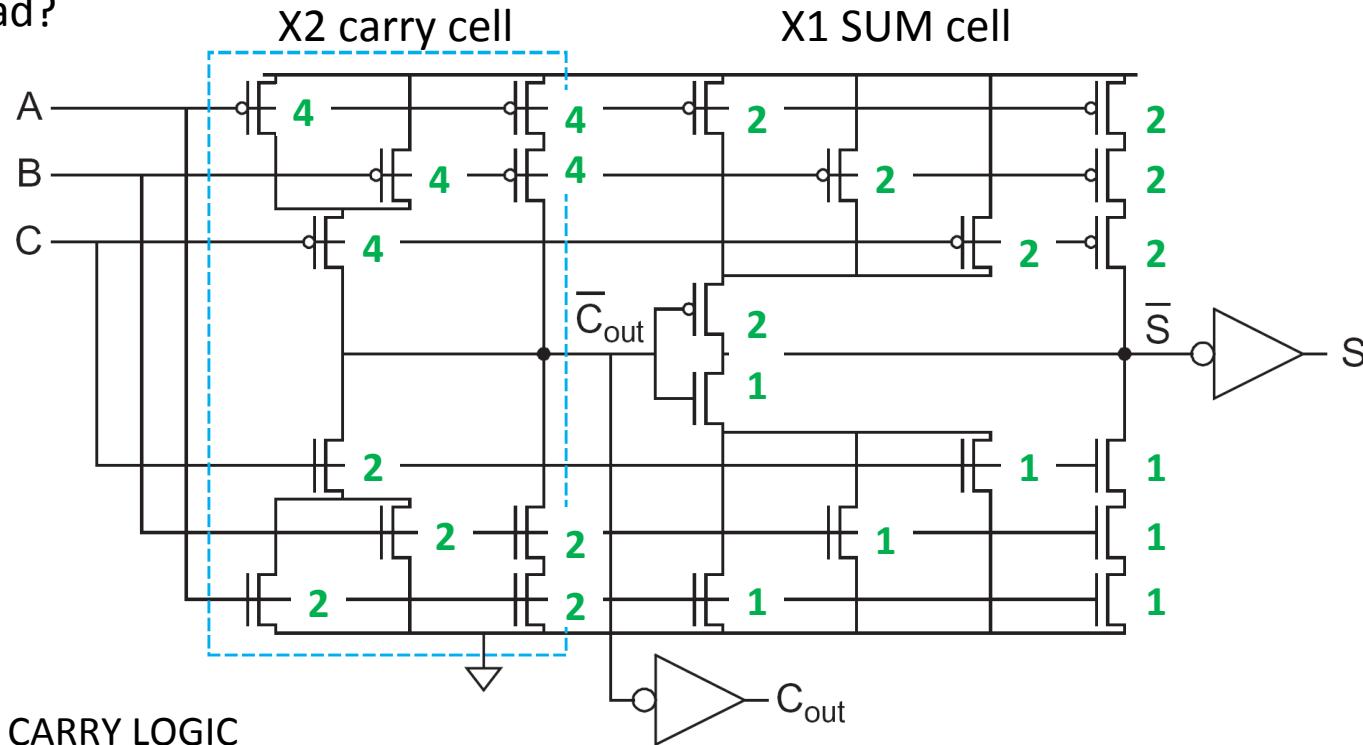
CARRY DELAY (CIN input)

$$d=4+2x(6+12)/6+1+(6+6+6)/12=4+6+1+1.5=12.5$$



Textbook solution

What if we resize
MOSFETs in SUM cell
for less load?



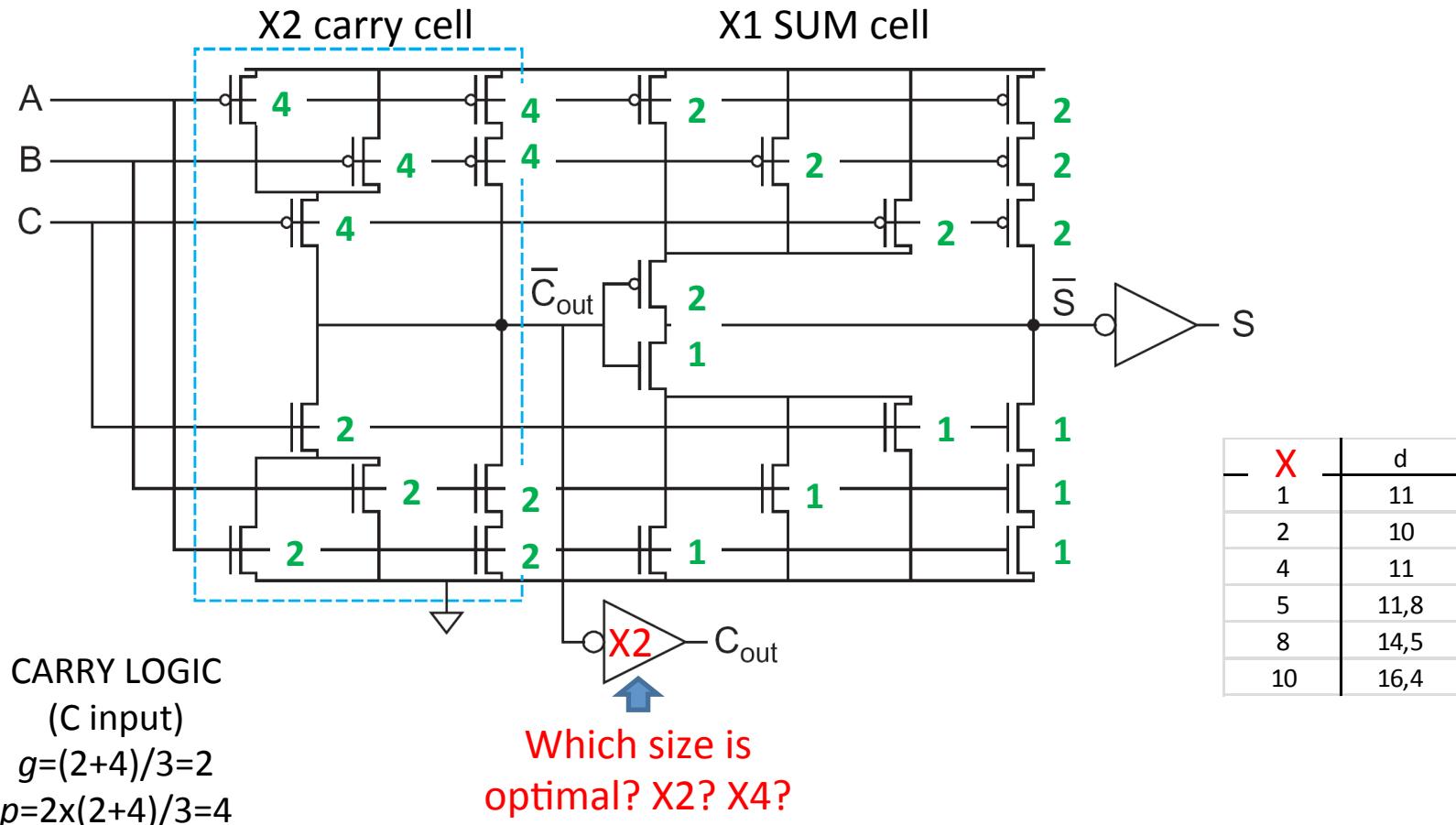
8-bit adder delay: Consider also SUM cell load

$$d = 8 \times (4 + 2 \times (1+2)/2 + 1 + (2+1+1)) = 8 \times (4 + 3 + 1) = 8 \times 10 = 80$$

$$\rightarrow t_{pd} = 400 \text{ ps}$$

Textbook solution

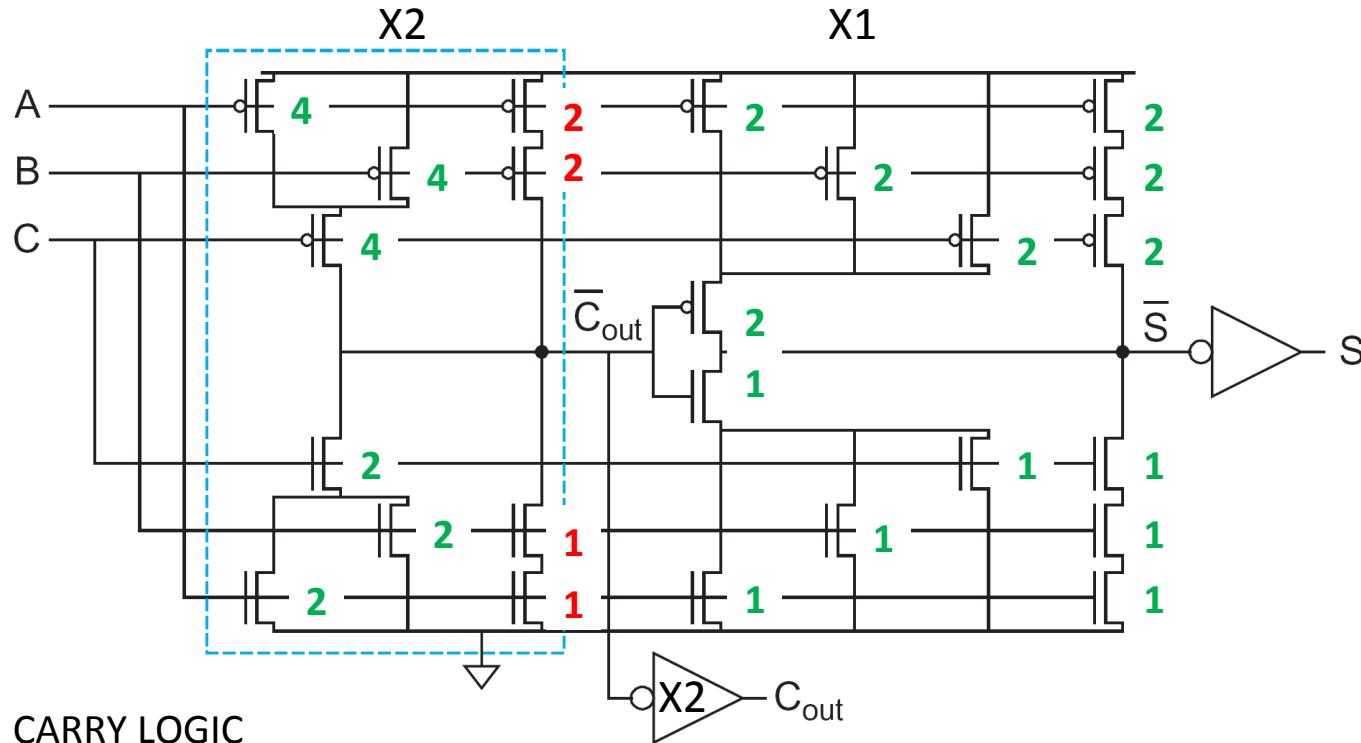
What if we resize
MOSFETs in SUM cell?



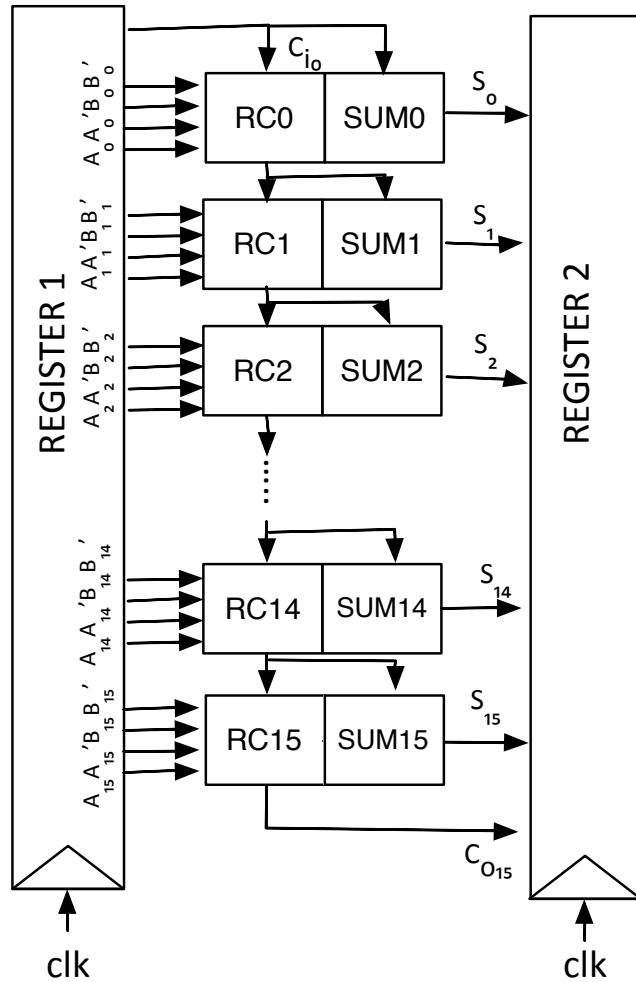
Textbook solution

CARRY DELAY (CIN input)

What if we also resize the MOSFETs as shown in red, what is the parasitic delay, p , now?



What is the critical path for entire adder?



What happens at
the end of the
adder?

Precomputing from A and B simplifies carry (and sum) computing

		AB				
		00	01	11	10	
Cin		0	0	1	0	Kill: $C_{out} = 0$
0	1	0	1	1	1	Generate: $C_{out} = 1$
1	0	1	1	1	1	Propagate: $C_{out} = C_{in}$
Kill		Propagate	Generate	Propagate		

Preview for week 6 & 7

Precomputing from A and B simplifies carry (and sum) computing

How to calculate P and G:

Generate:

$$C_{out} = I$$

Propagate:

$$C_{out} = C_{in}$$

Generate: $G = AB$

Propagate: $P = A \oplus B$

Can be done for all bits immediately

How to use P and G in full adder:

$$C_{out} = G + PC_{in}$$

$$Sum = P \oplus C_{in}$$

Which operation has the longest propagation delay?

Preview for week 6& 7

Conclusion

- The test vectors matter.
 - What is the worst case?
- Optimal sizing of inverter helps
 - Removing inverter entirely would be even better
- Global optimization of critical path
 - Size down circuits not in critical path to limit their capacitive load
- Precomputing “intermediate signals” from A and B also helps

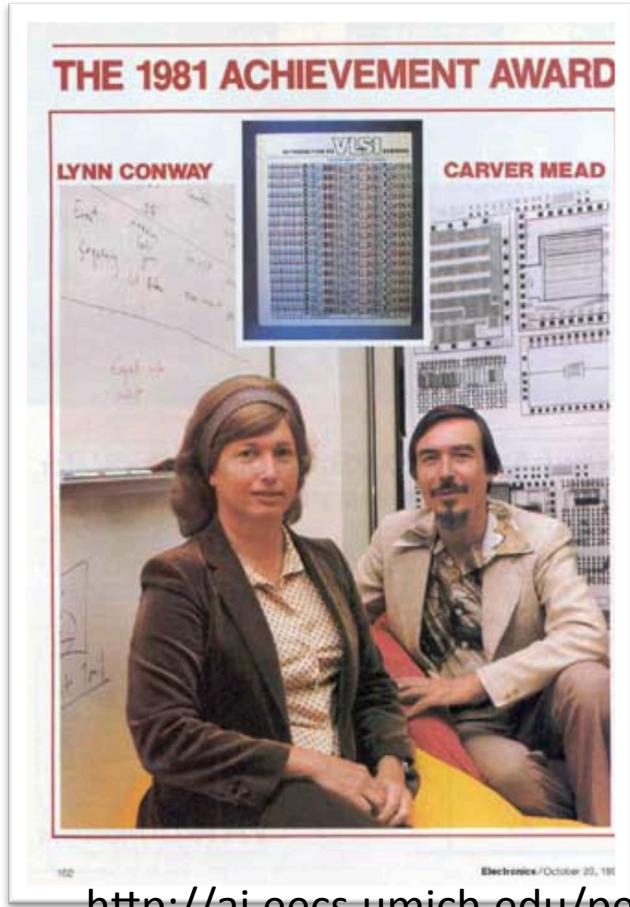
Aim of the lecture

- To give some basic understanding of intralayer and interlayer geometric design rules
 - Intralayer rules concerns minimum dimensions and spacing between objects in **the same** layer
 - Interlayer rules concerns minimum spacing and overlap rules between objects in two **different** layers:
 - MOSFET rules between poly and active
 - Contact and via rules between contacts or vias to their bottom and top contacting layers
- Discuss the prelab assignment for the layout lab session

Why rules?

- To minimize the risk of faulty circuits after fabrication.
 - Malfunction due to shorts
 - Improper forming of components or contacts
 - Electrical charging during fabrication.
- More complex nowadays due to:
 - Smaller feature sizes
 - More complex fabrication processes
 - For example polishing for each metal layer

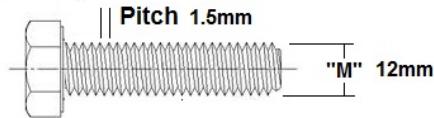
Who invented the simple (lambda) CMOS rules?



Carver Mead
Lynn Conway

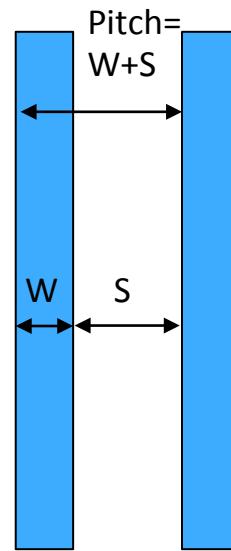
http://ai.eecs.umich.edu/people/conway/Memoirs/VLSI/SSCM/VLSI_Reminiscences.pdf

Geometric design rules



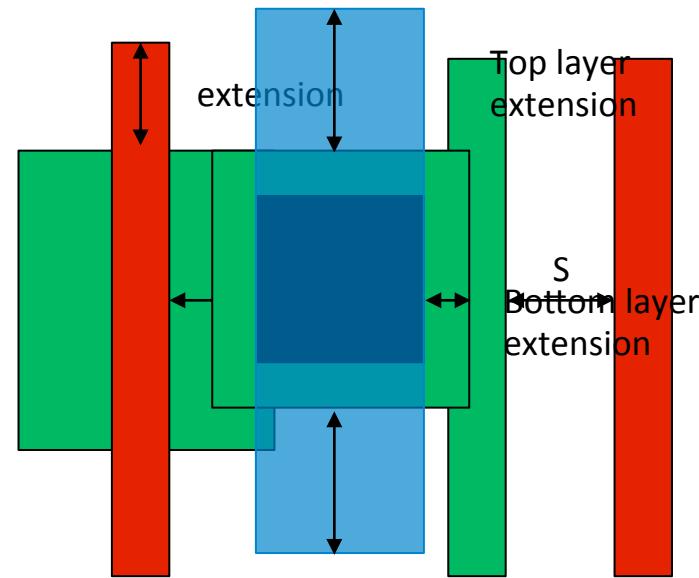
Bolt identification
"M" Dia*Pitch

M12*1.5



Interlayer rules

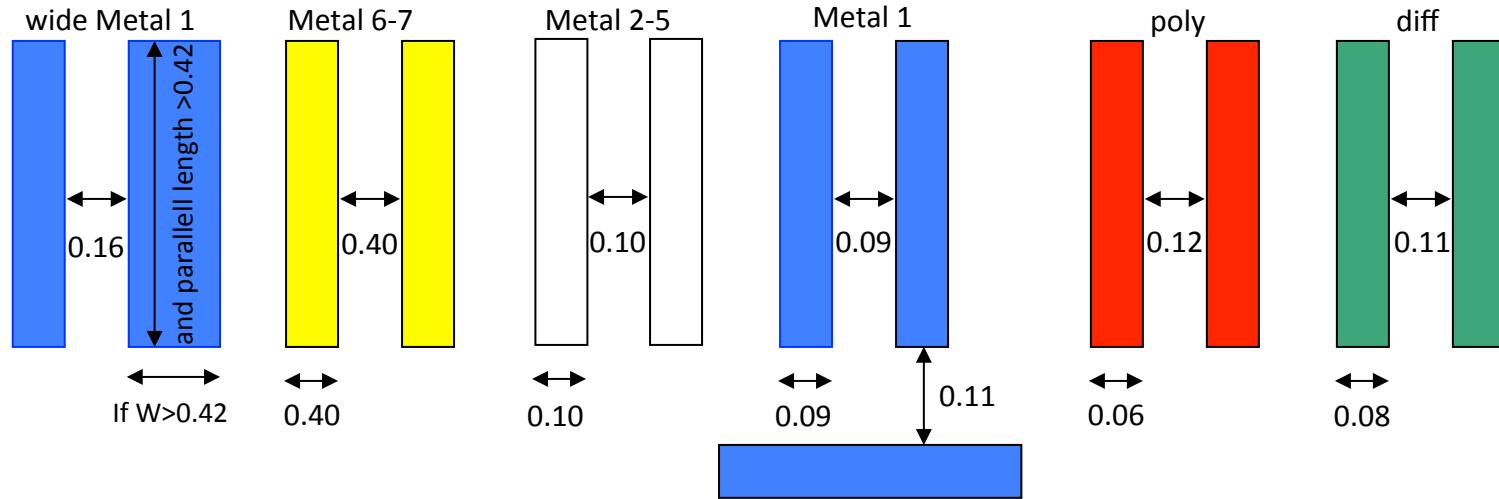
Anti-tack polyvias



pitch = width + separation

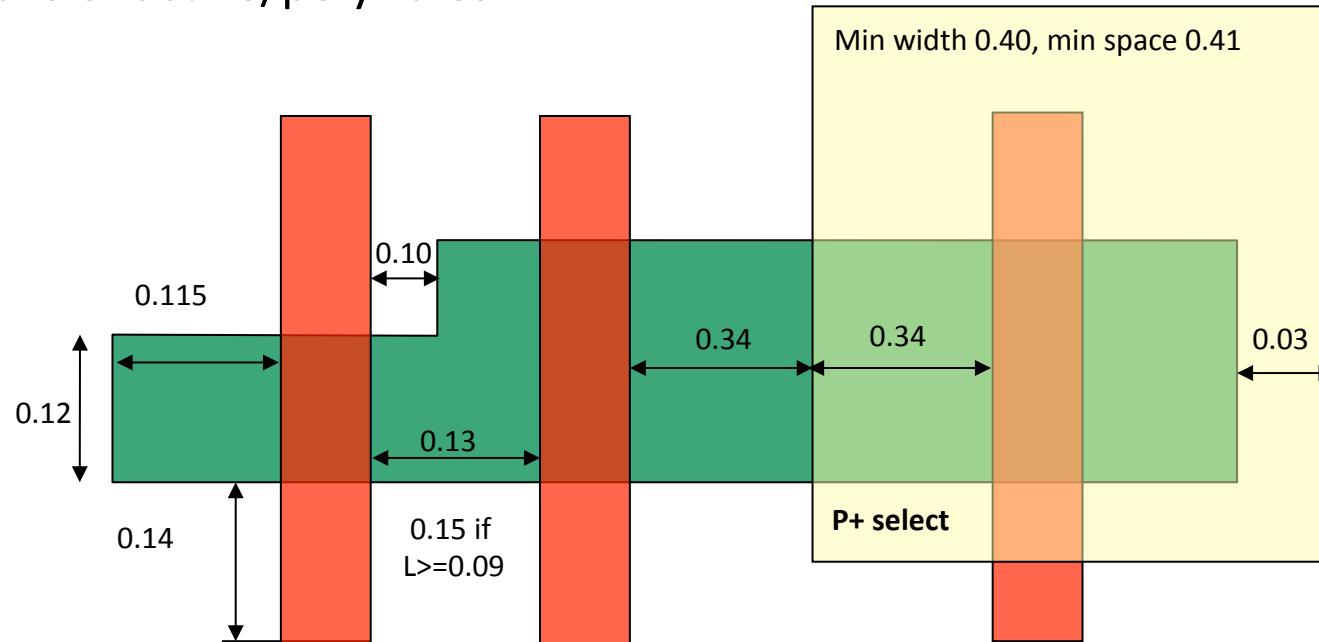
Geometric design rules

Intralayer rules



Geometric design rules

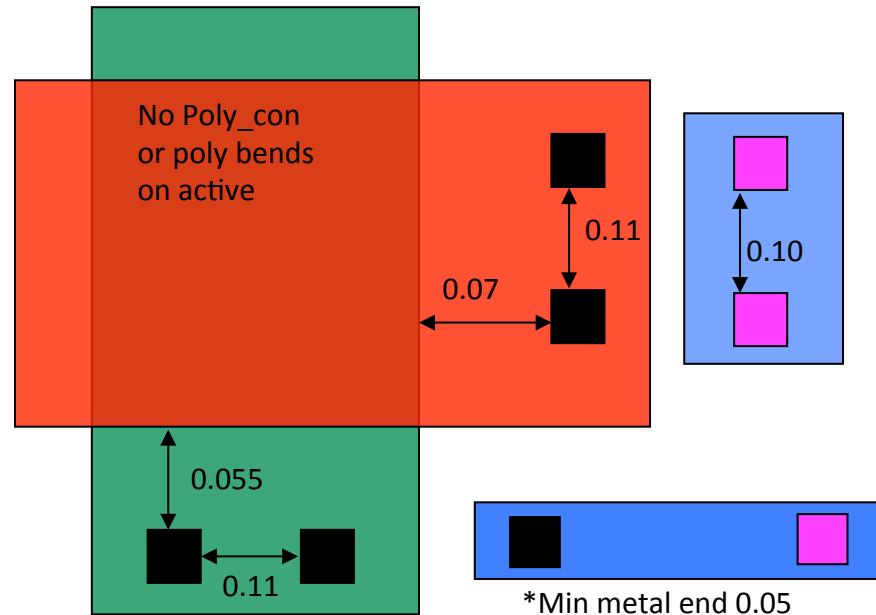
Details of active/poly rules



Geometric design rules

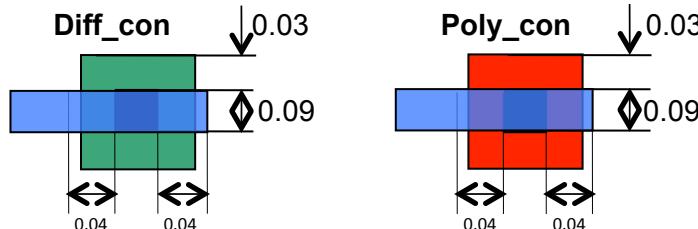
Contact and via* rules

*A via is a contact between two metal layers.



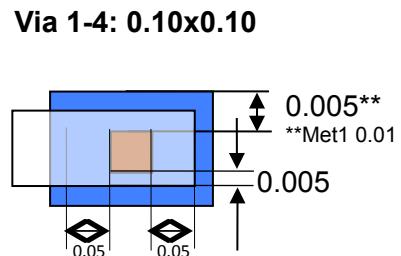
Geometric design rules

Contact and via rules

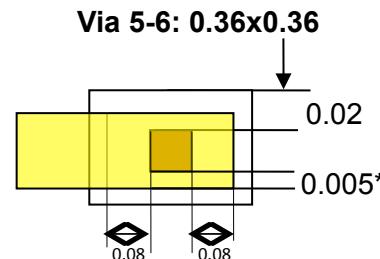


Also allowed to enclose contact with metal1 0.025 on all four sides.

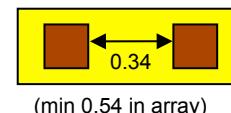
For minimum width metal lines at least 0.04 enclosure is required on two opposite sides (bottom measures).



At least two 0.05 metal ends on metal pad

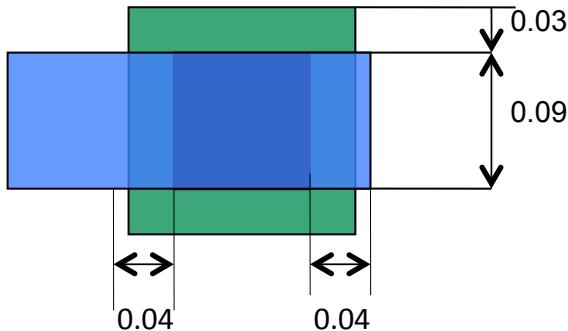


At least two 0.08 metal ends on metal pad

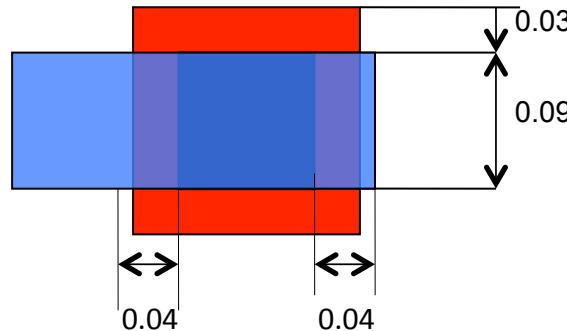


Detailed contacts

Diffusion contact



Poly contact



The contacts drawn here are the ones we recommend.

For minimum width metal lines (width = 0.09) at least 0.04 enclosure is required on two opposite sides, while on remaining two sides enclosure is 0.0.

Options:

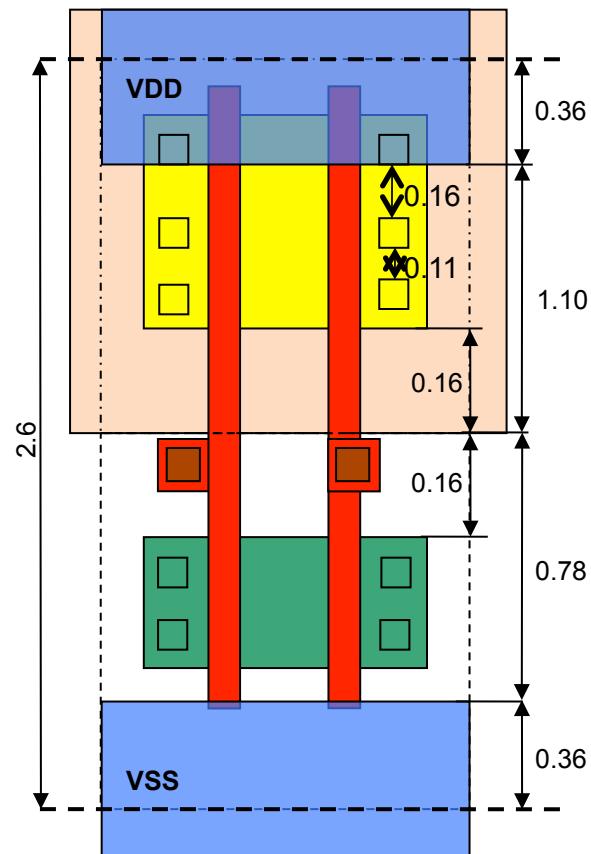
It is also allowed to have metal1 enclose contact 0.025 in all four directions.

Diffusion rectangular: extension has to be 0.03 on two opposite sides and 0.015 on the two remaining sides.

Poly rectangular: extension has to be 0.04 on two opposite sides and 0.01 on the two remaining sides.

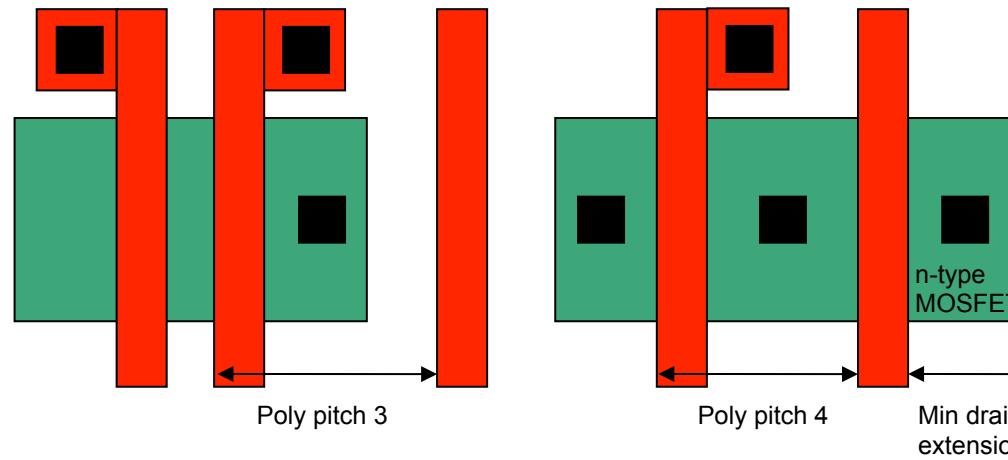
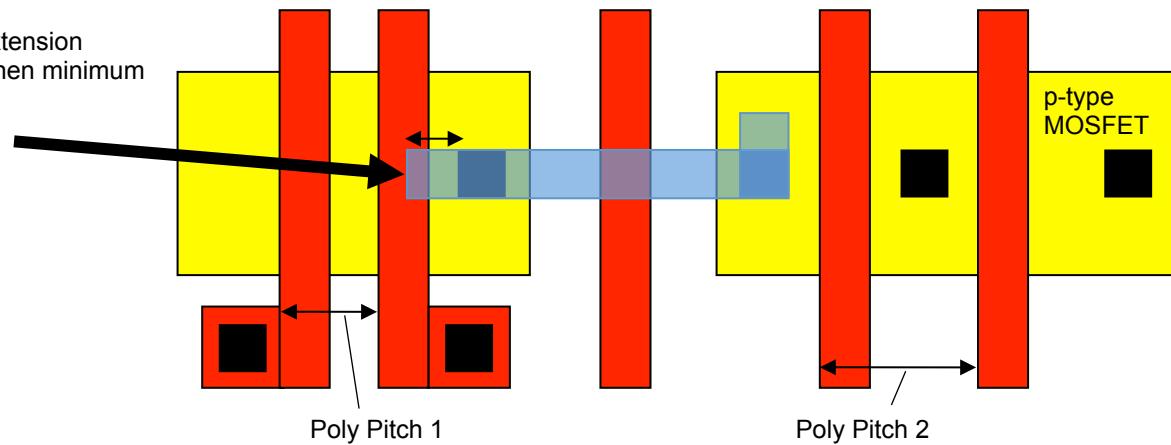
Geometric design rules

Standard cell template
with rails, n-well, pplus and nplus design rules

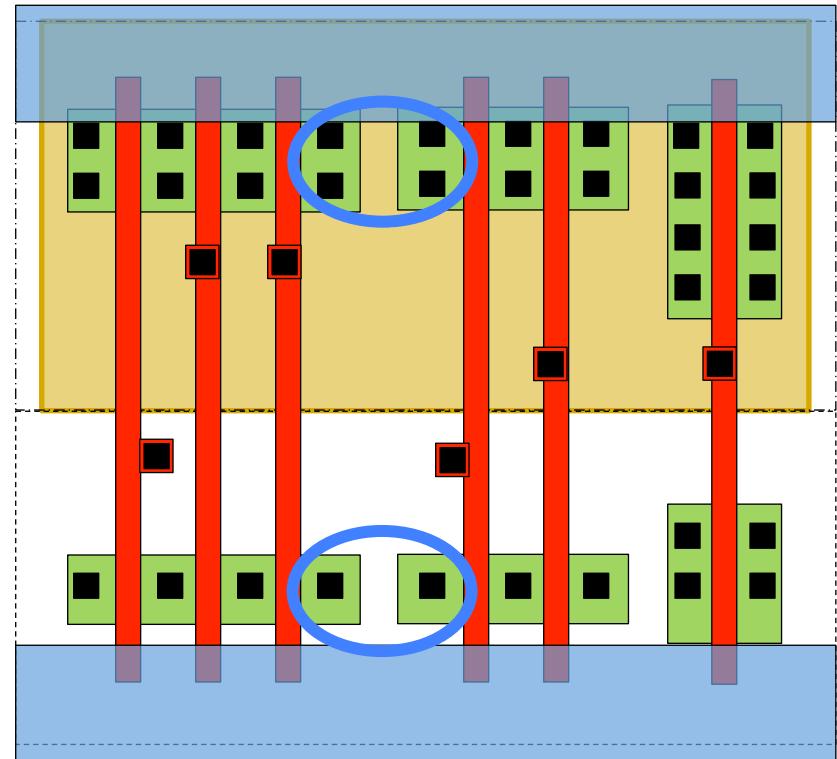
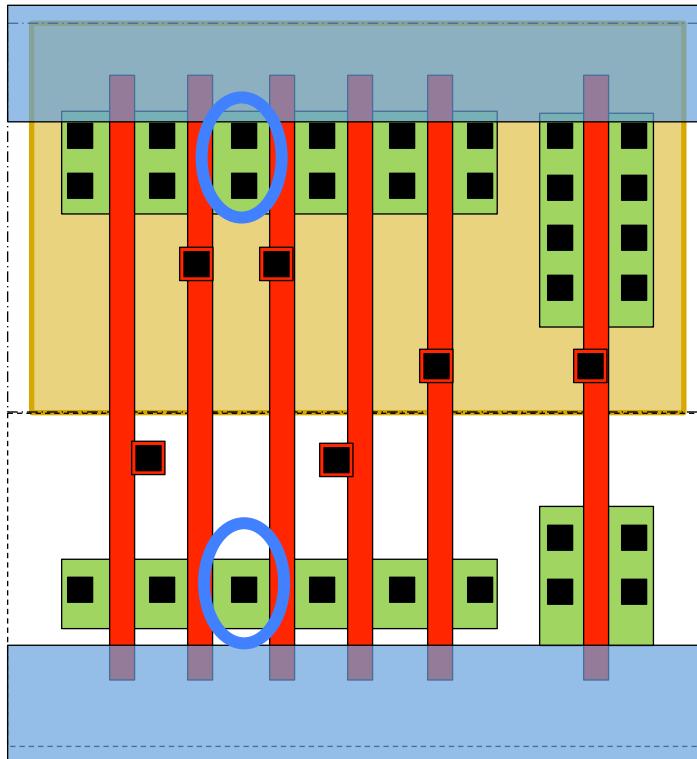


Prelab 3 assignment

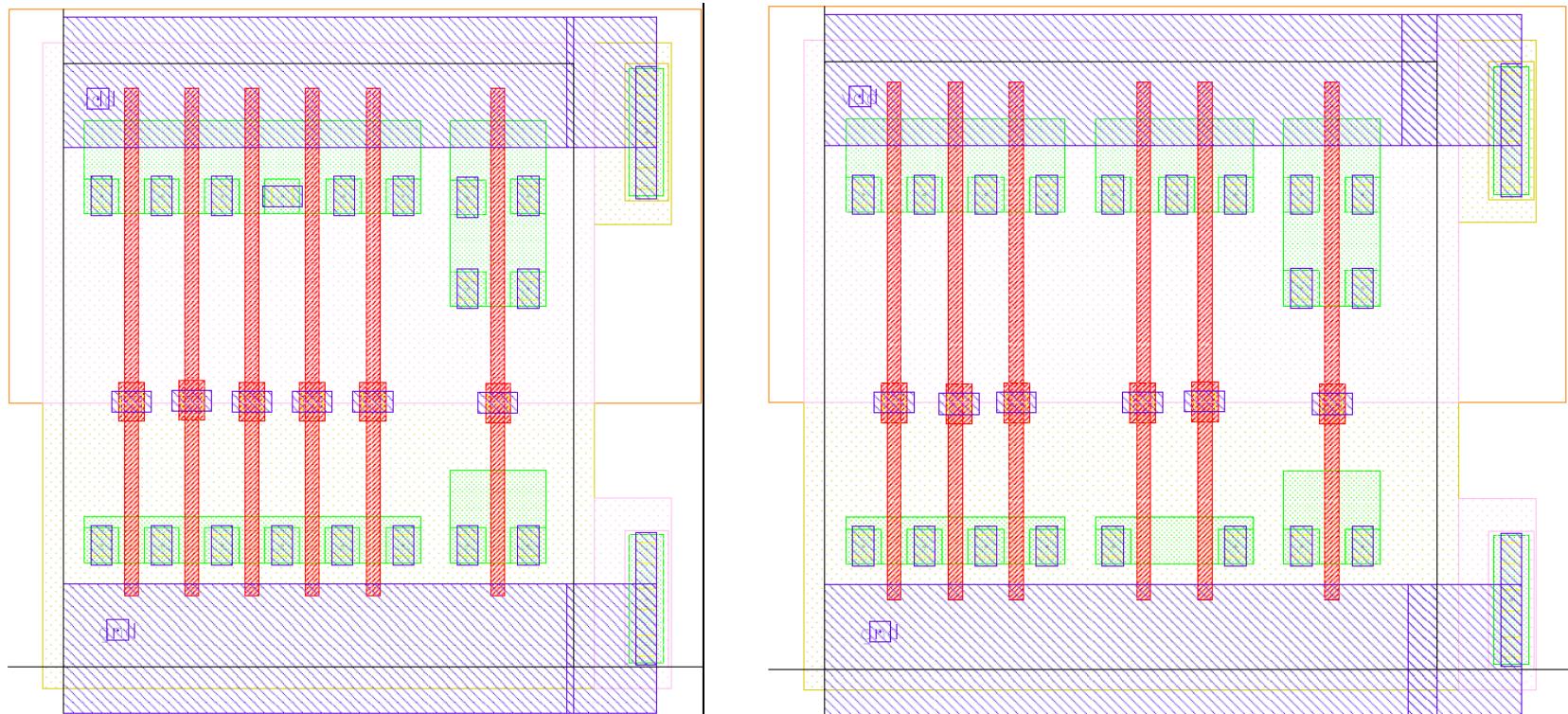
Minimum metal extension beyond contact when minimum wire width is used



Prelab 3 layout templates



Prelab 3 layout templates

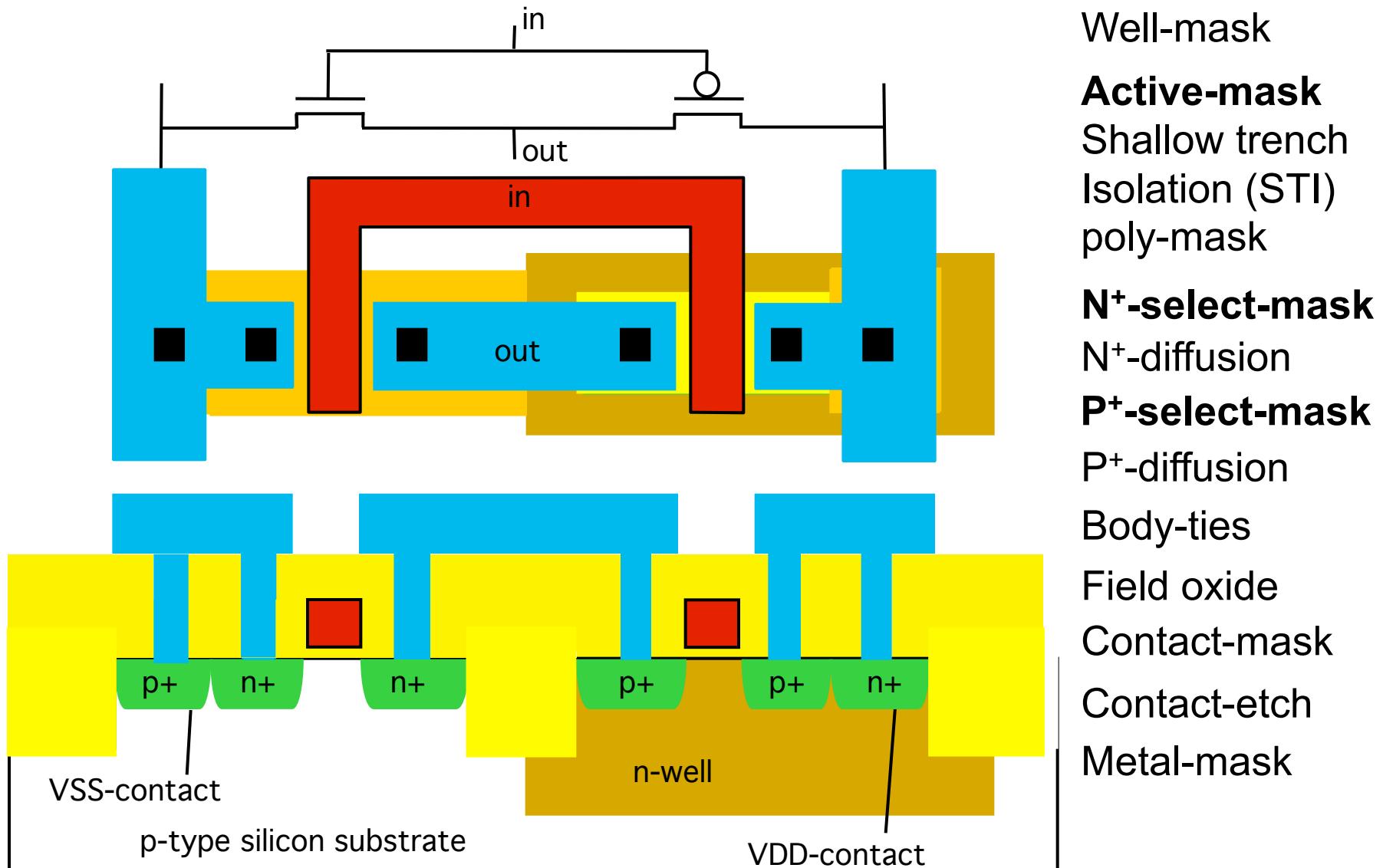


The layers

- **Physical layers** in your layout:
 - N-well (NW)
 - Diffusion/Active (OD)
 - Poly (PO)
 - **Metal-1 (M1)**
 - Contact hole (CO)
- **Non-physical layers** required in your layout:
 - P-select (PPLUS) (pink dots)
 - N-select (NPLUS) (yellow dots)

Example of layers

CMOS Layout



Summary

In this lecture, we have

- introduced intralayer and interlayer geometric design rules
- defined the minimum pitch:
 - the sum of the minimum width and the minimum separation between two objects in the **same** layer
- studied interlayer rules between objects in two different layers, like
 - MOSFET rules between poly and active
 - contact and via rules between contacts or vias to their bottom and top contacting layers
- discussed the prelab assignment to the layout lab session